



Politecnico  
di Torino



e-Lite



# Introduction to Graphs

Tecniche di Programmazione – A.A. 2021/2022





# Definition: Graph

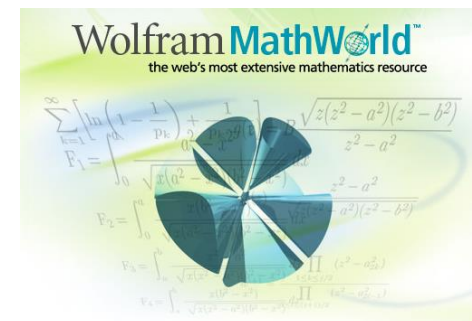
Introduction to Graphs

# Definition: **Graph**

---

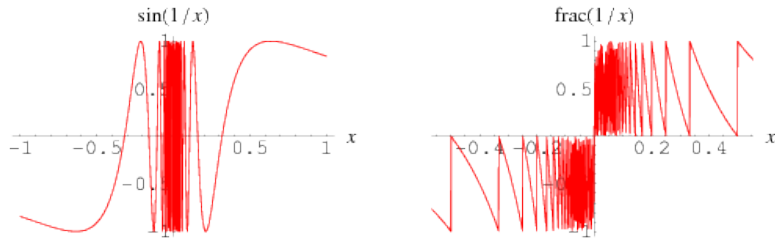
- ▶ A **graph** is a collection of **points** and **lines** connecting some (possibly empty) subset of them.
- ▶ The points of a graph are most commonly known as **graph vertices**, but may also be called “nodes” or simply “points.”
- ▶ The lines connecting the vertices of a graph are most commonly known as **graph edges**, but may also be called “arcs” or “lines.”

<http://mathworld.wolfram.com/>

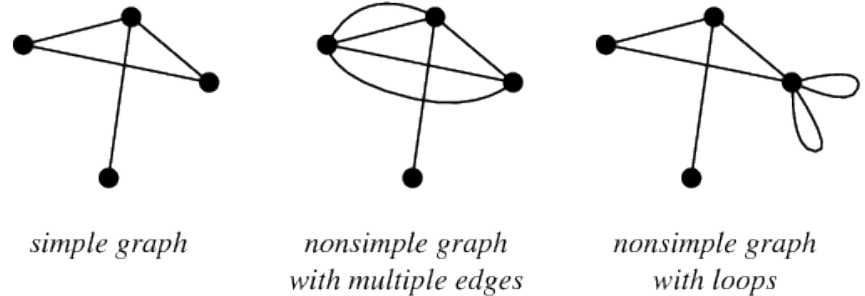


# Big warning: Graph $\neq$ Graph $\neq$ Graph

**Graph (plot)**  
(italiano: grafico)

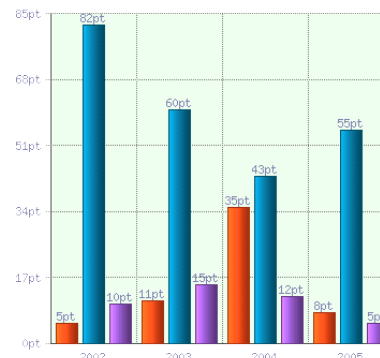


**Graph (maths)**  
(italiano: grafo)



$\neq$

**Graph (chart)**  
(italiano: grafico)



# History

---

- ▶ The study of graphs is known as **graph theory**, and was first systematically investigated by D. König in the 1930s
- ▶ Euler's proof about the *walk across all seven bridges of Königsberg* (1736), now known as the *Königsberg bridge problem*, is a famous precursor to graph theory.
- ▶ In fact, the study of various sorts of paths in graphs has many applications in real-world problems.

# Königsberg Bridge Problem

- ▶ Can the 7 bridges of the city of Königsberg over the river Pregel all be traversed in a single trip without doubling back, with the additional requirement that the trip ends in the same place it began?

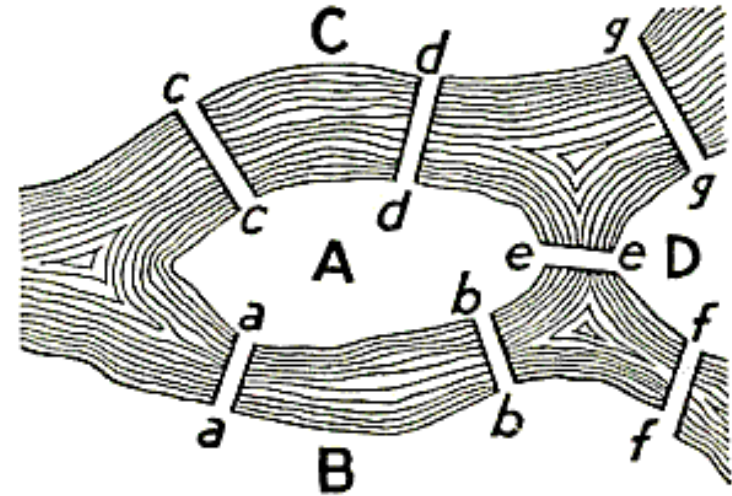
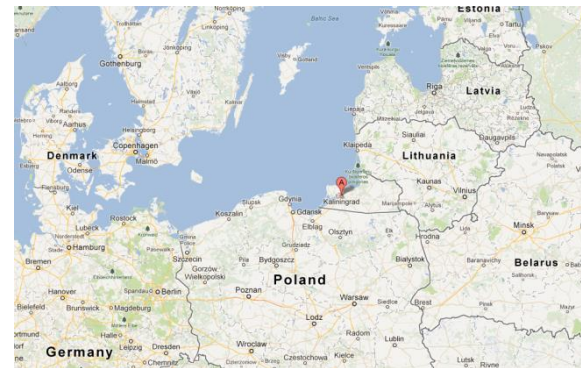


FIGURE 98. *Geographic Map:  
The Königsberg Bridges.*



Today: Kaliningrad, Russia

# Königsberg Bridge Problem

- ▶ Can the 7 bridges of the city of Königsberg over the river Pregel all be traversed in a single trip without crossing any bridge twice, and returning to the starting point?

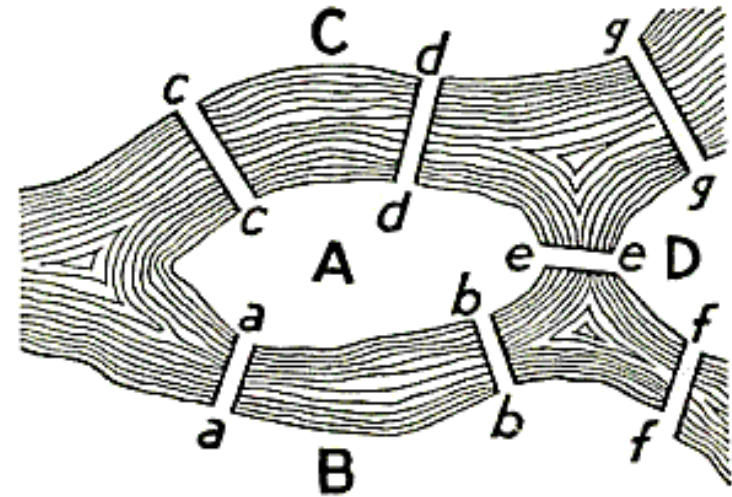
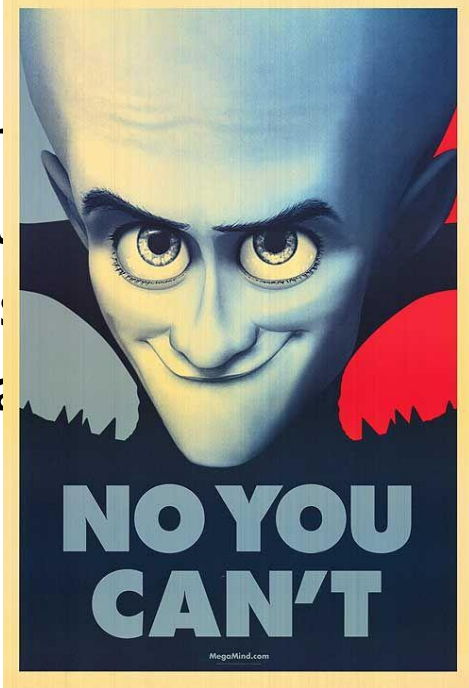
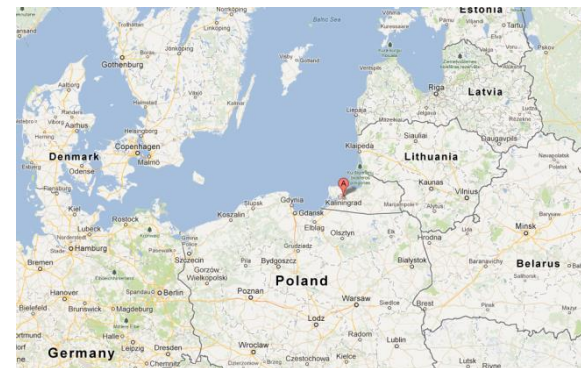


FIGURE 98. *Geographic Map:  
The Königsberg Bridges.*



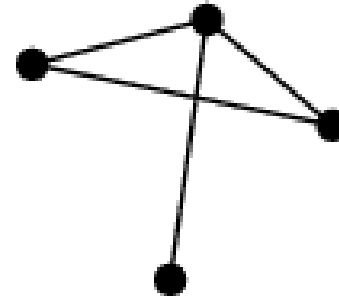
Today: Kaliningrad, Russia

# Types of graphs: edge cardinality

---

- ▶ **Simple graph:**

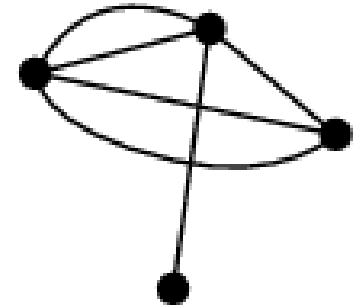
- ▶ At most one edge (i.e., either one edge or no edges) may connect any two vertices



*simple graph*

- ▶ **Multigraph:**

- ▶ Multiple edges are allowed between vertices



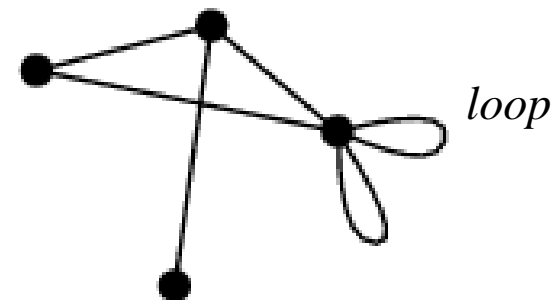
*multigraph*

- ▶ **Loops:**

- ▶ Edge between a vertex and itself

- ▶ **Pseudograph:**

- ▶ Multigraph with loops

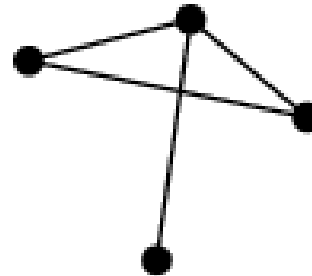


*pseudograph*

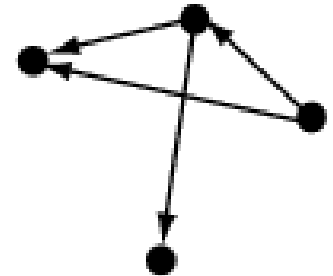


# Types of graphs: edge direction

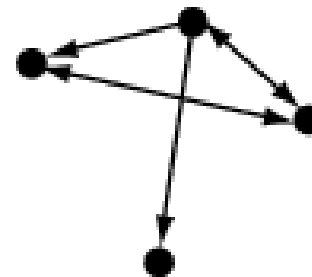
- ▶ Undirected
- ▶ Oriented
  - ▶ Edges have **one** direction (indicated by arrow)
- ▶ Directed
  - ▶ Edges may have **one or two** directions
- ▶ Network
  - ▶ Oriented graph with weighted edges



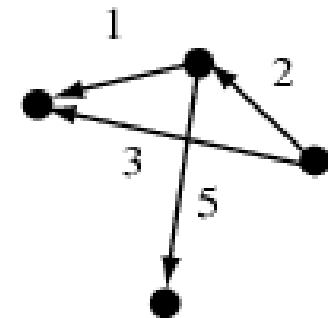
*undirected graph*



*oriented graph*



*directed graph*

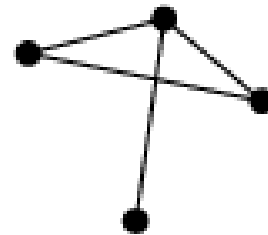


*network*

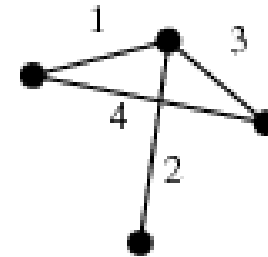
# Types of graphs: labeling

## ▶ Labels

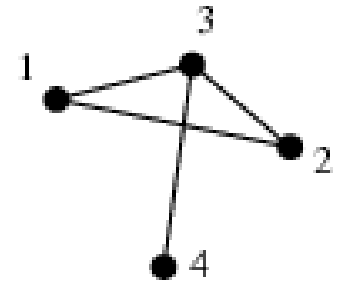
- ▶ None
- ▶ On Vertices
- ▶ On Edges



*unlabeled graph*



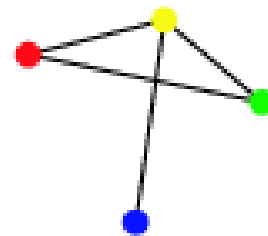
*edge-labeled graph*



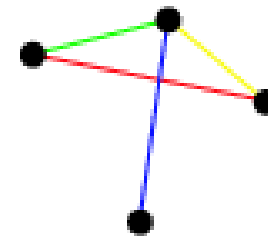
*vertex-labeled graph*

## ▶ Groups (=colors)

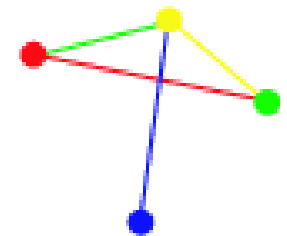
- ▶ Of Vertices
  - ▶ no edge connects two identically colored vertices
- ▶ Of Edges
  - ▶ adjacent edges must receive different colors
- ▶ Of both



*vertex-colored graph*



*edge-colored graph*

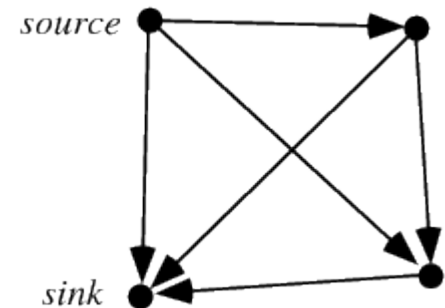


*vertex- and edge-colored graph*

# Directed and Oriented graphs

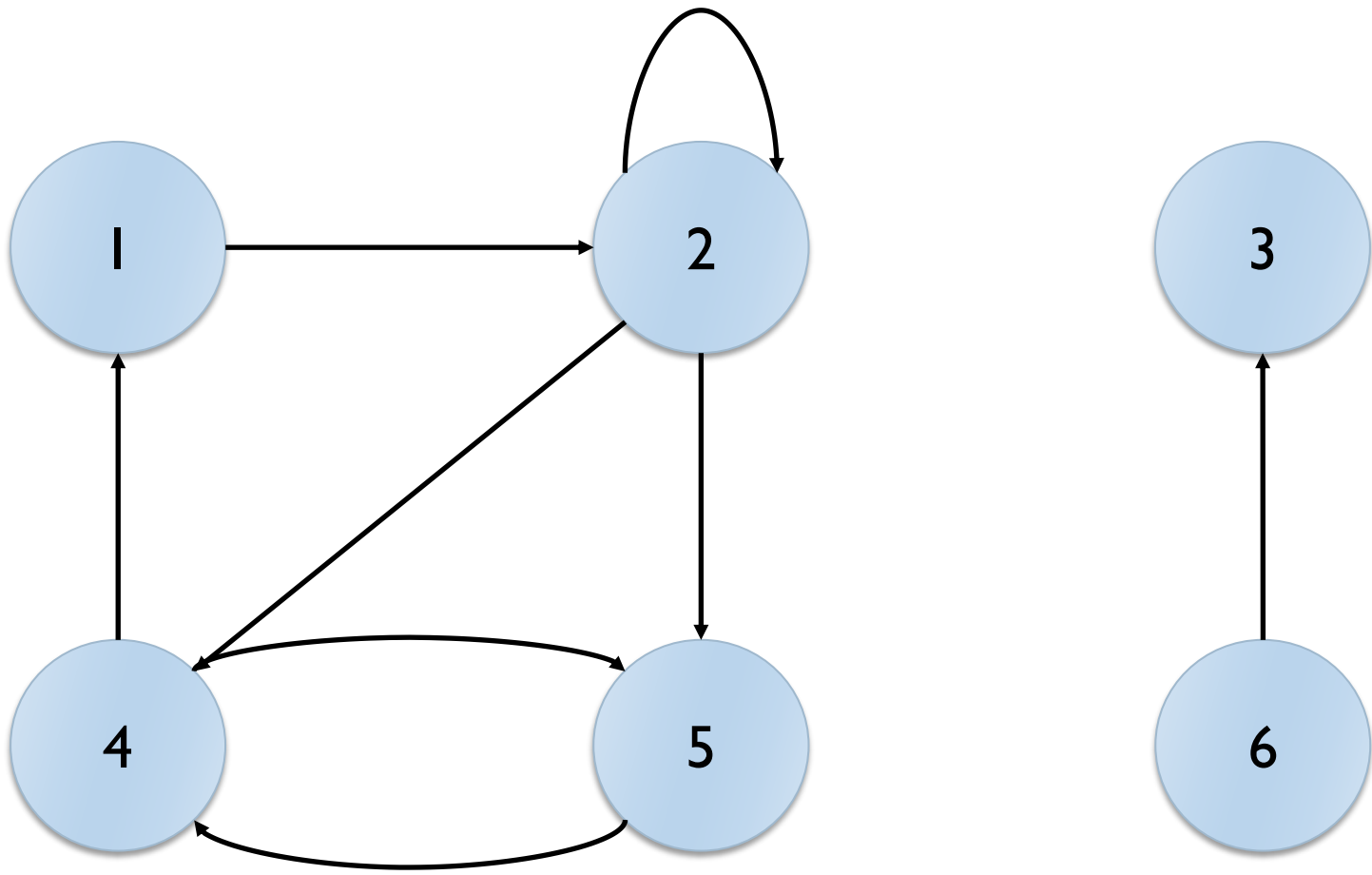
---

- ▶ A Directed Graph (*di-graph*)  $G$  is a pair  $(V,E)$ , where
  - ▶  $V$  is a (finite) set of *vertices*
  - ▶  $E$  is a (finite) set of *edges*, that identify a binary relationship over  $V$ 
    - ▶  $E \subseteq V \times V$



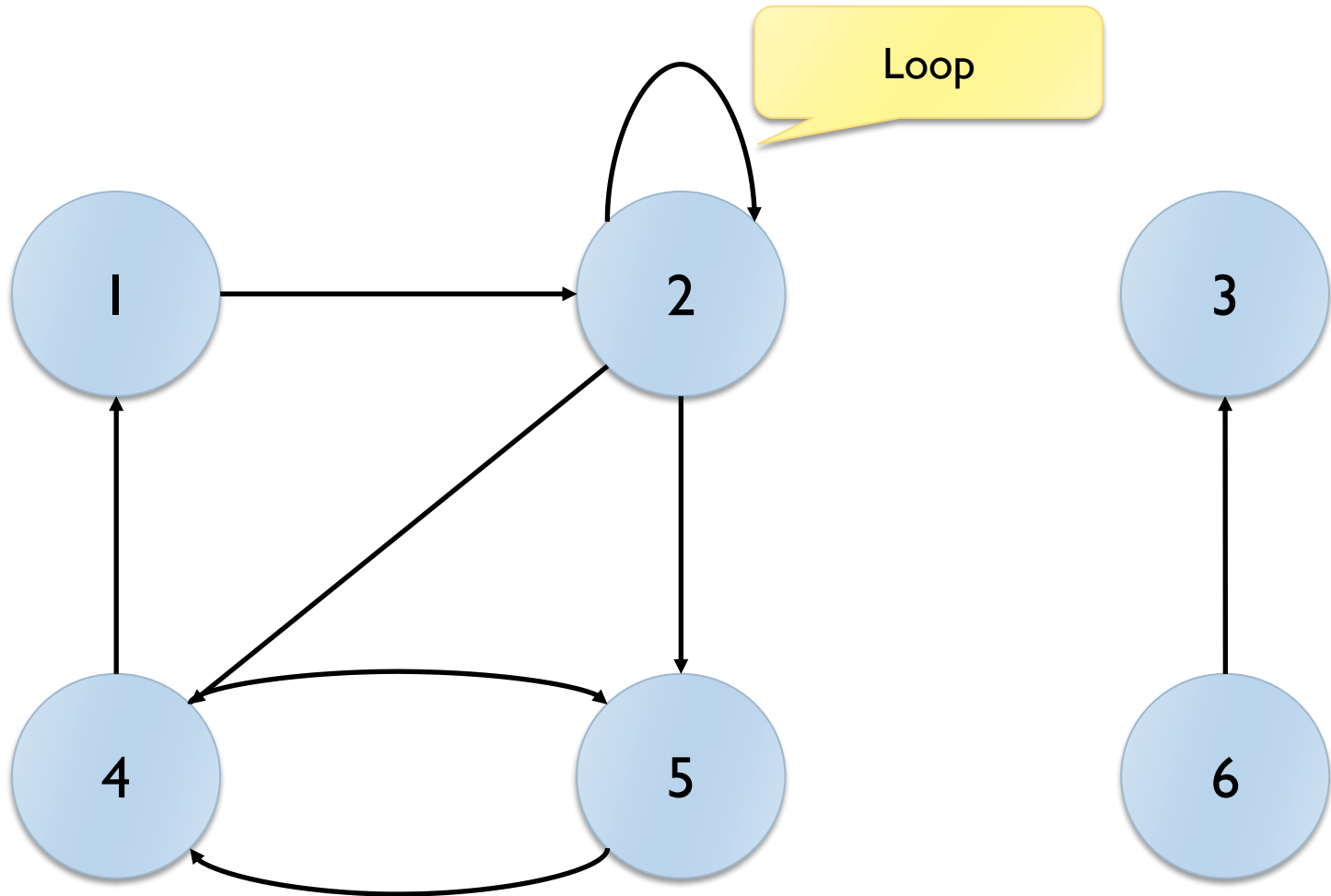
# Example

---



# Example

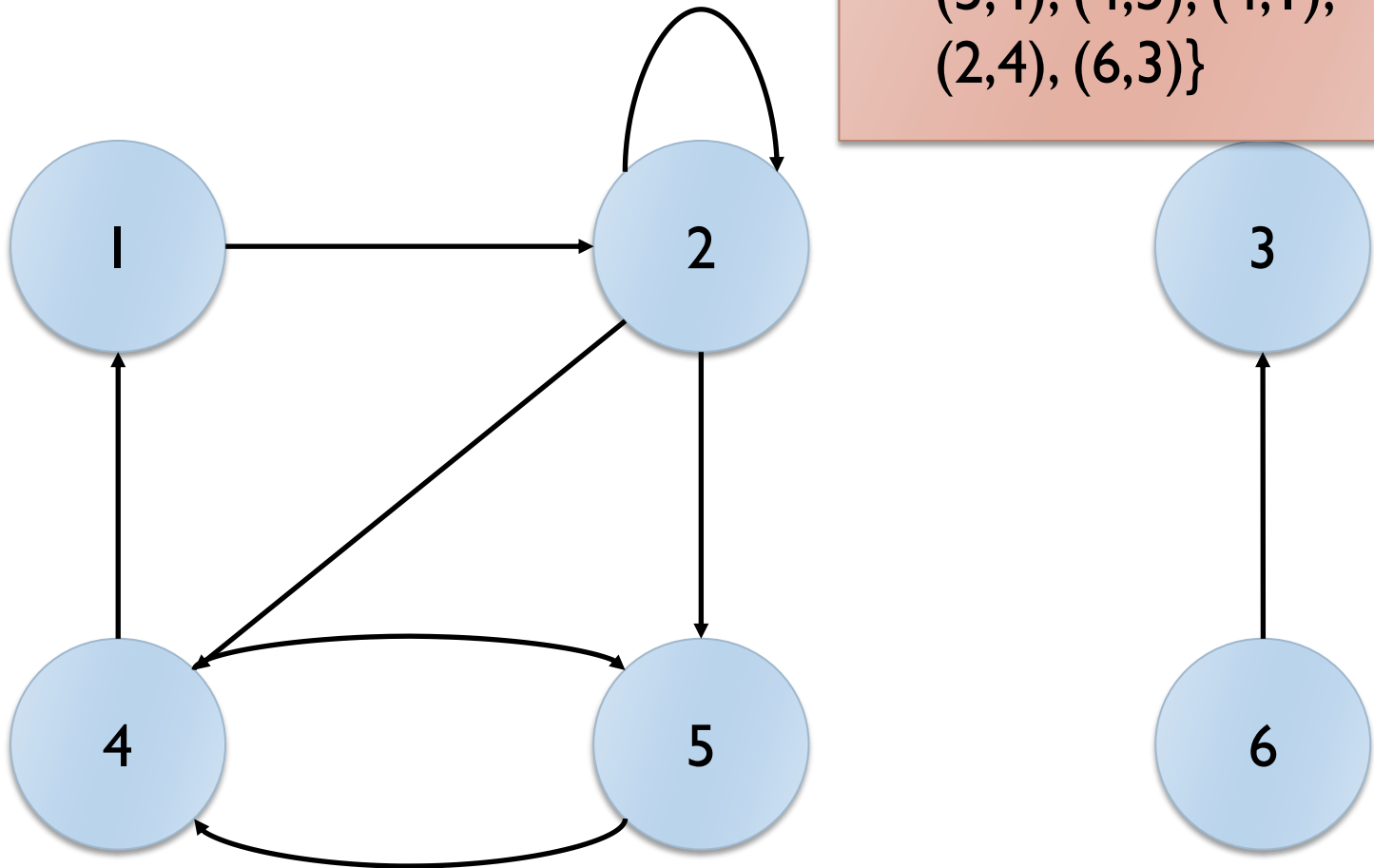
---



# Example

$$V = \{1, 2, 3, 4, 5, 6\}$$

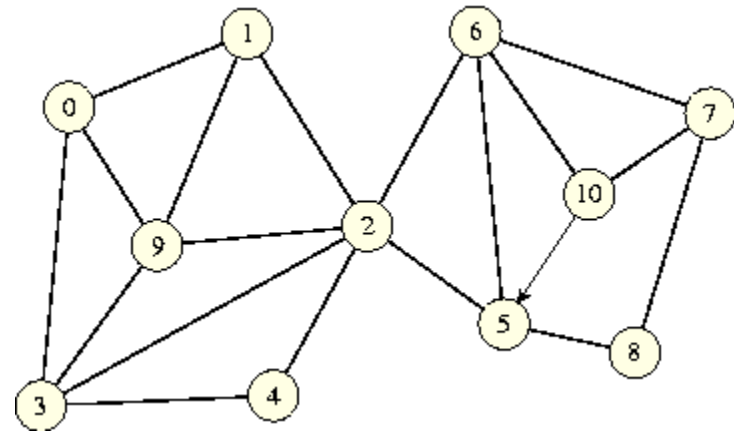
$$E = \{(1, 2), (2, 2), (2, 5), (5, 4), (4, 5), (4, 1), (2, 4), (6, 3)\}$$



# Undirected graph

---

- ▶ An **Undirected** Graph is still represented as a couple  $G=(V,E)$ , but the set  $E$  is made of **non-ordered pairs** of vertices

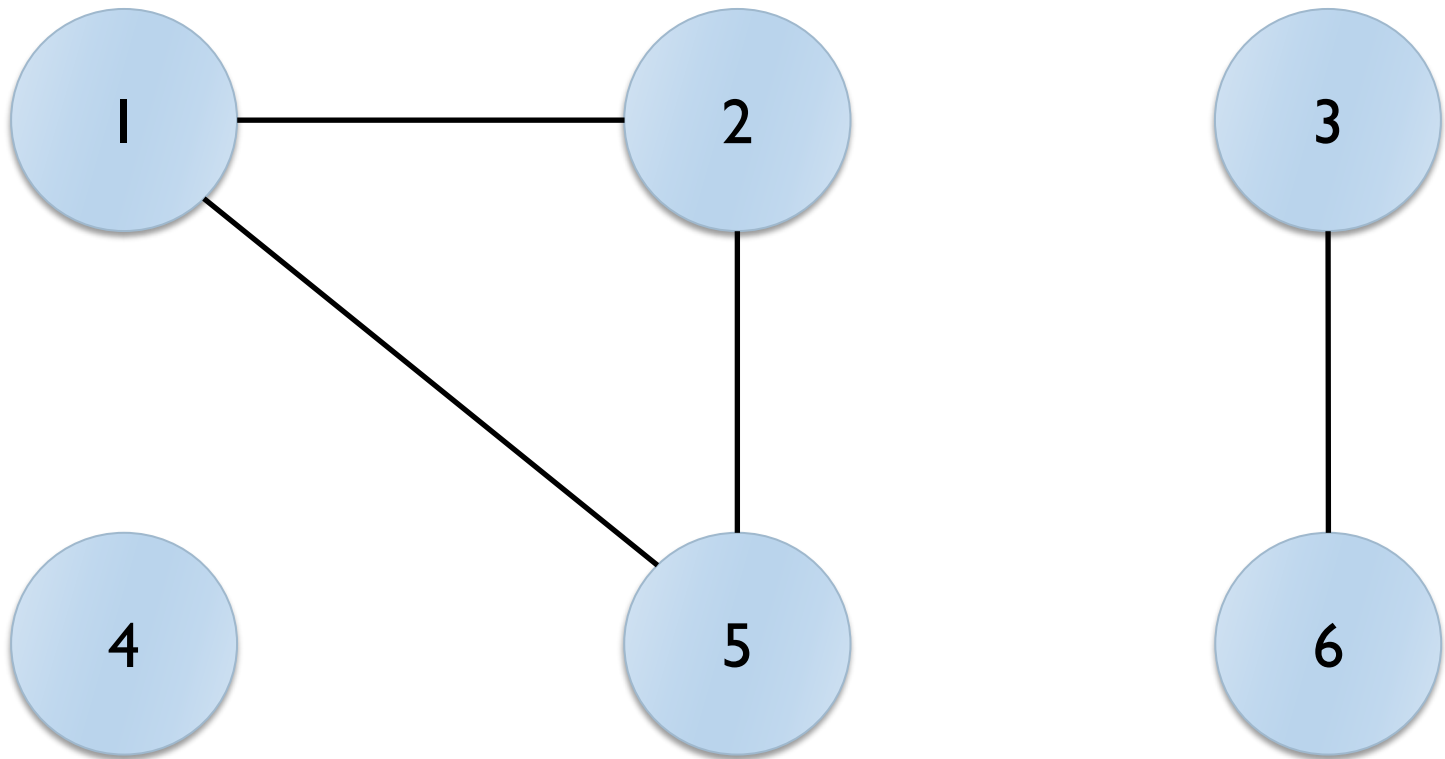


# Example

---

$V = \{1, 2, 3, 4, 5, 6\}$

$E = \{\{1, 2\}, \{2, 5\}, \{5, 1\}, \{6, 3\}\}$

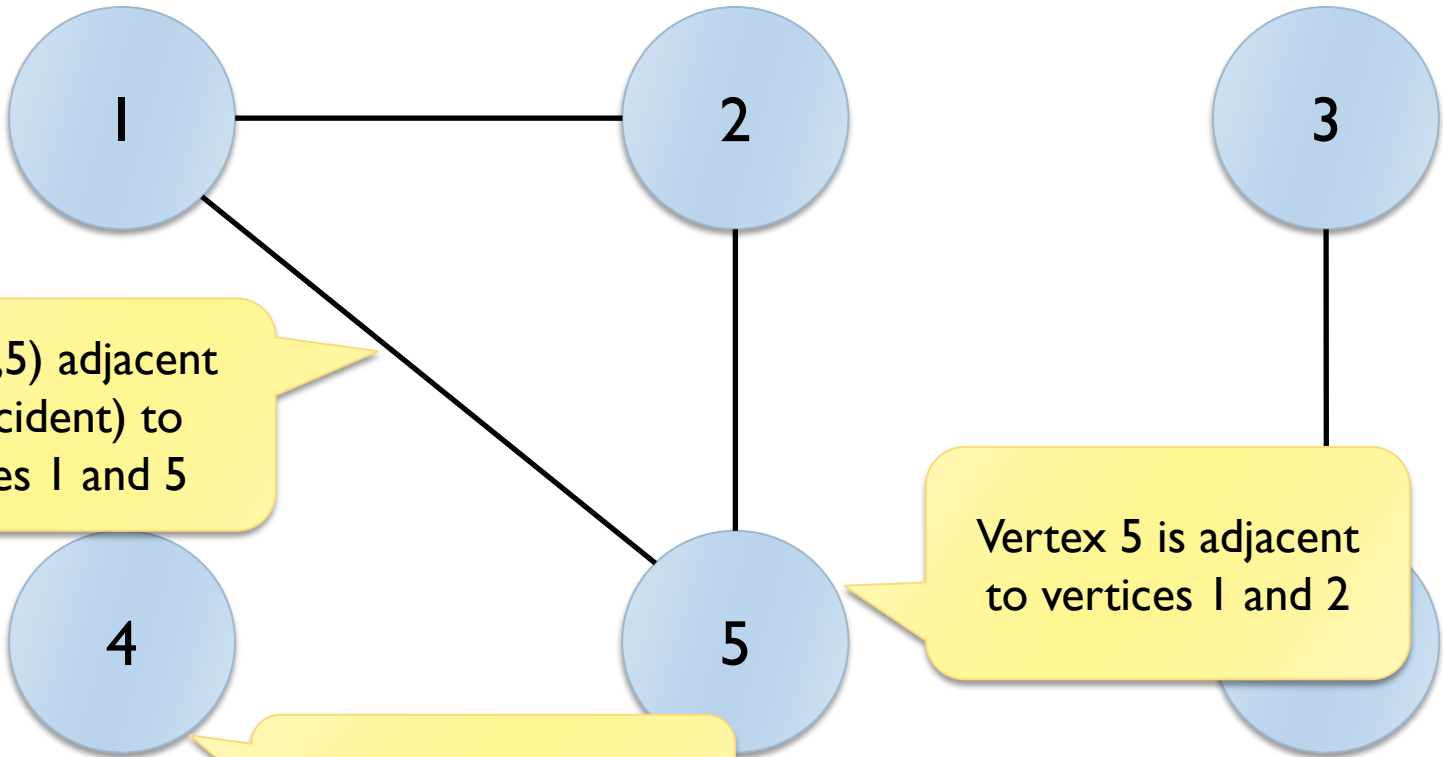




# Example

$V = \{1, 2, 3, 4, 5, 6\}$

$E = \{\{1, 2\}, \{2, 5\}, \{5, 1\}, \{6, 3\}\}$



Edge (1,5) adjacent  
(or incident) to  
vertices 1 and 5

Vertex 5 is adjacent  
to vertices 1 and 2

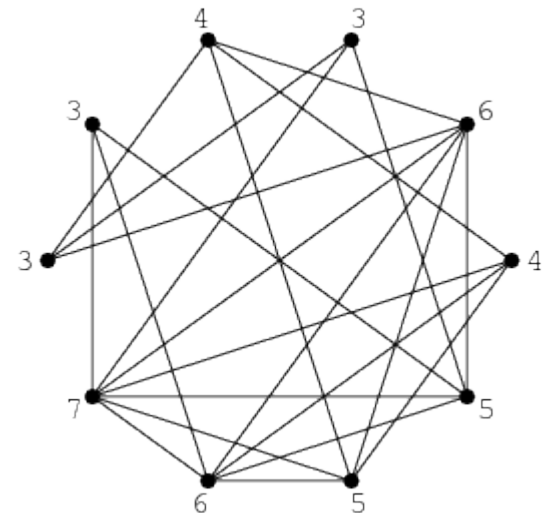
Vertex 4 is isolated



# Degree

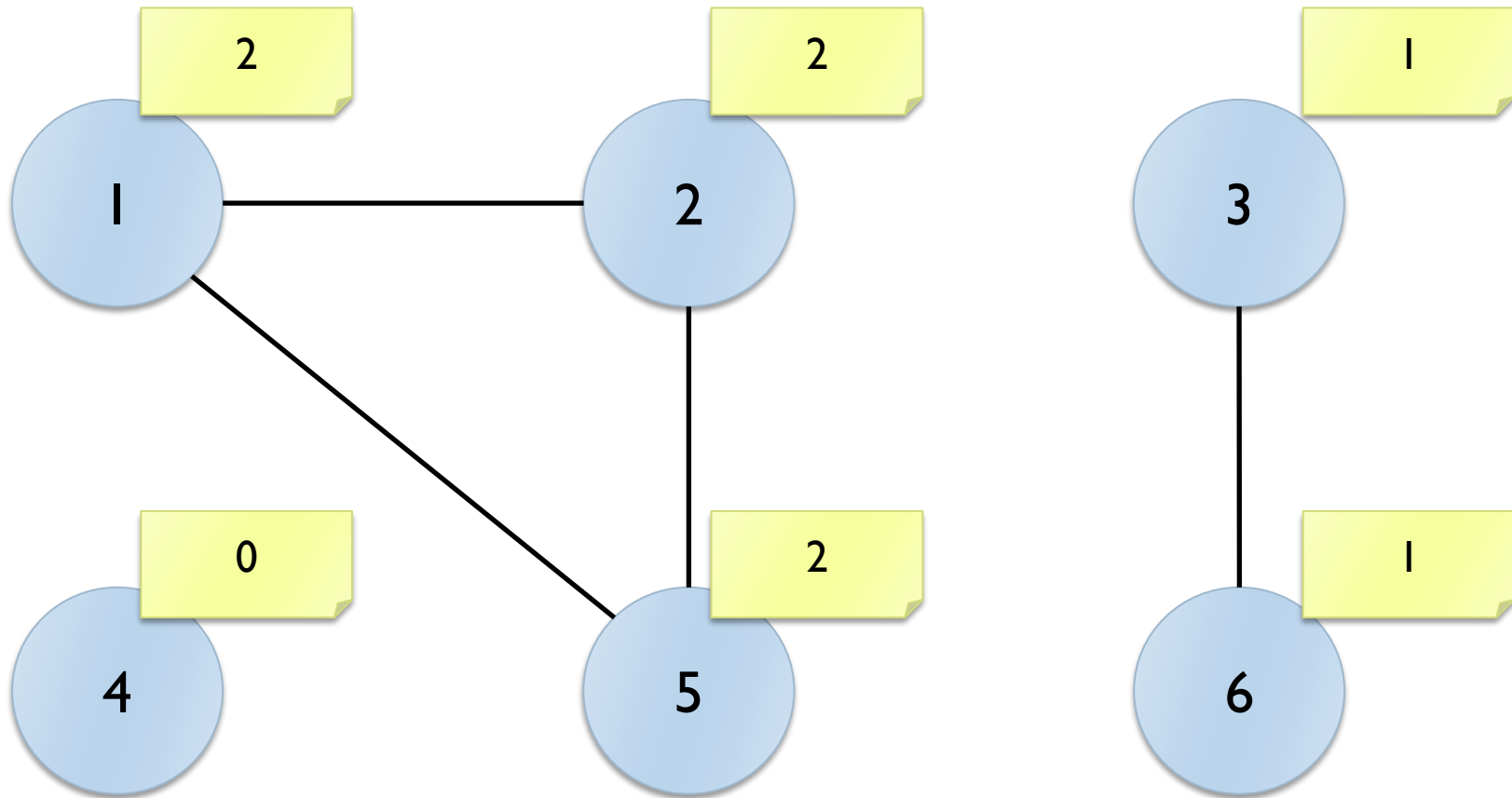
---

- ▶ In an *undirected* graph,
  - ▶ the **degree** of a vertex is the number of incident edges
- ▶ In a *directed* graph
  - ▶ The **in-degree** is the number of incoming edges
  - ▶ The **out-degree** is the number of departing edges
  - ▶ The **degree** is the sum of in-degree and out-degree
- ▶ A vertex with degree 0 is **isolated**

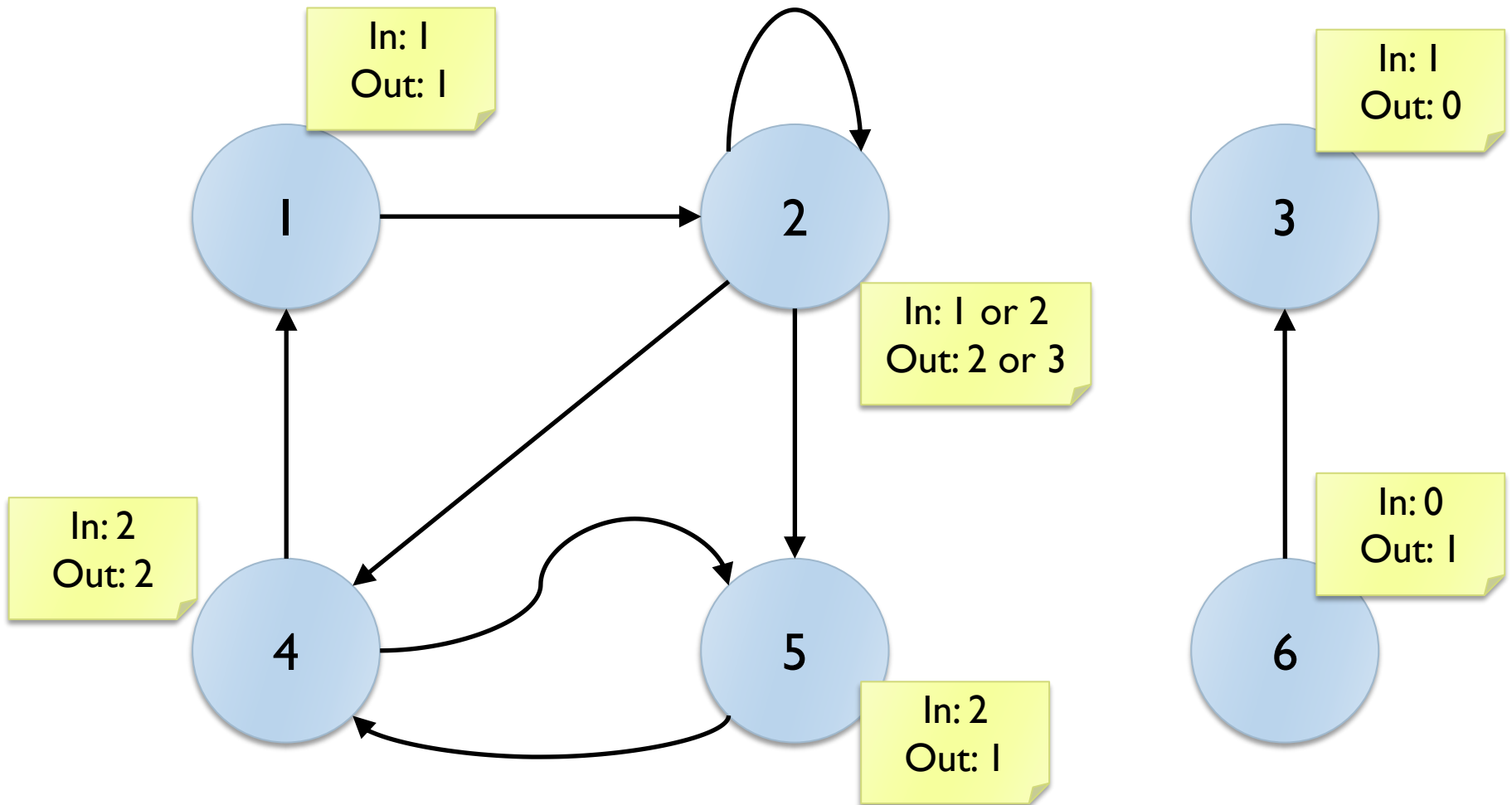


# Degree

---



# Degree



# Paths

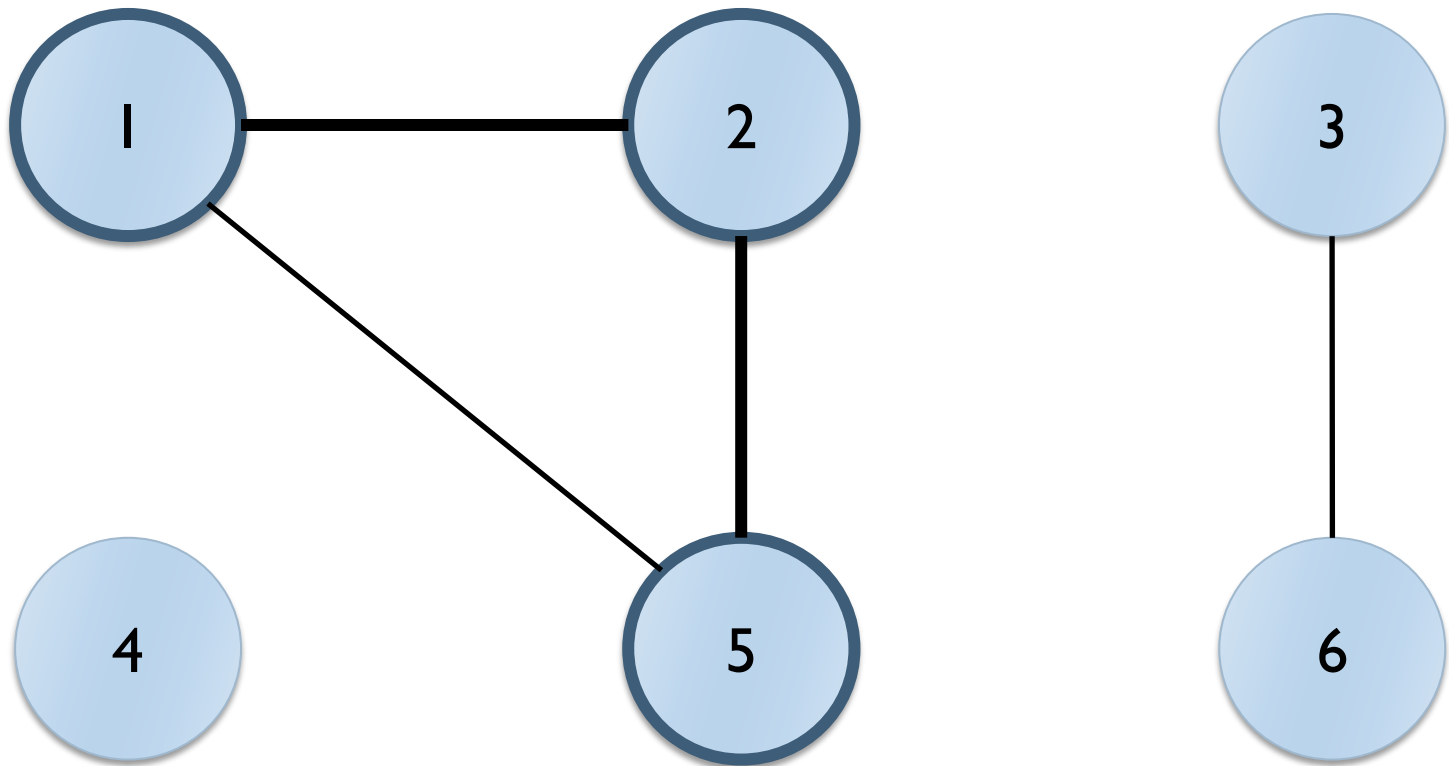
---

- ▶ A **path** on a graph  $G=(V,E)$  also called a trail, is a sequence  $\{v_1, v_2, \dots, v_n\}$  such that:
  - ▶  $v_1, \dots, v_n$  are vertices:  $v_i \in V$
  - ▶  $(v_1, v_2), (v_2, v_3), \dots, (v_{n-1}, v_n)$  are graph edges:  $(v_{i-1}, v_i) \in E$
  - ▶  $v_i$  are distinct (for “simple” paths).
- ▶ The **length** of a path is the number of edges  $(n-1)$
- ▶ If there exist a path between  $v_A$  and  $v_B$  we say that  $v_B$  is **reachable** from  $v_A$

# Example

---

Path = { 1, 2, 5 }  
Length = 2



# Cycles

---

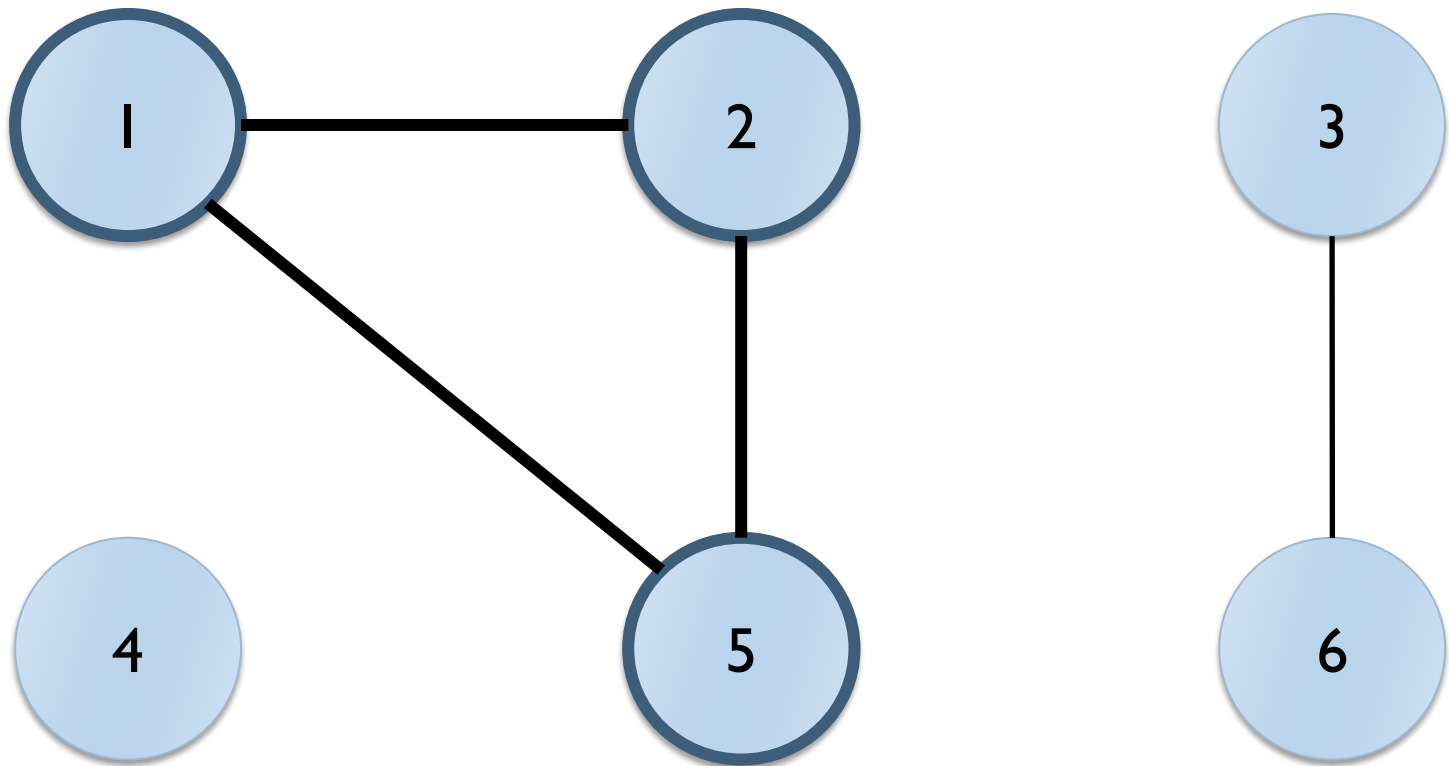
- ▶ A cycle is a path where  $v_1 = v_n$
- ▶ A graph with no cycles is said acyclic



# Example

---

Path = { 1, 2, 5, 1 }  
Length = 3



# Reachability (Undirected)

---

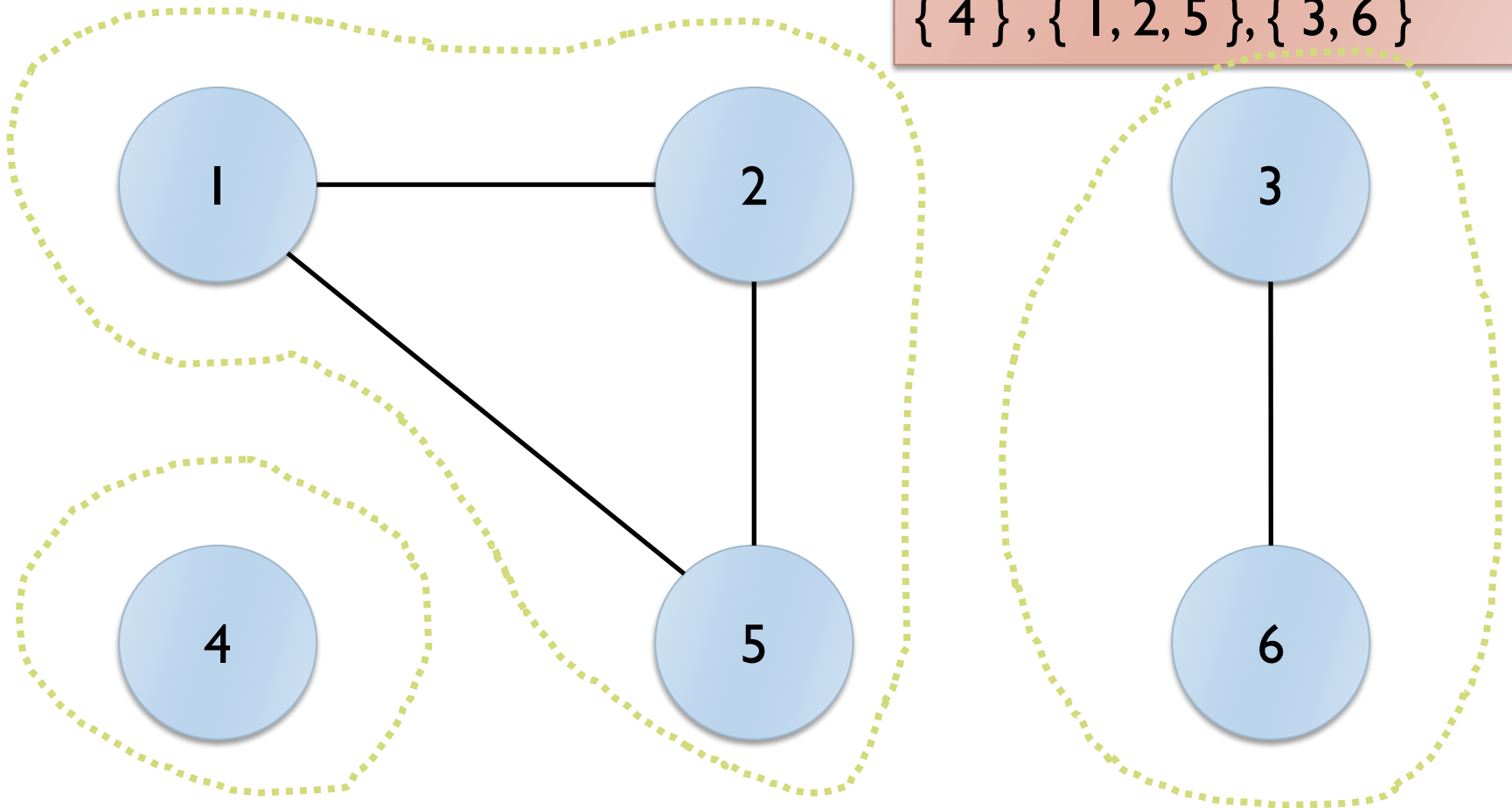
- ▶ An undirected graph is **connected** if, for every couple of vertices, there is a path connecting them
- ▶ The connected sub-graphs of maximum size are called **connected components**
- ▶ A connected graph has exactly one connected component

# Connected components

The graph is **not** connected.

Connected components =  
3

$\{ 4 \}, \{ 1, 2, 5 \}, \{ 3, 6 \}$



# Reachability (Directed)

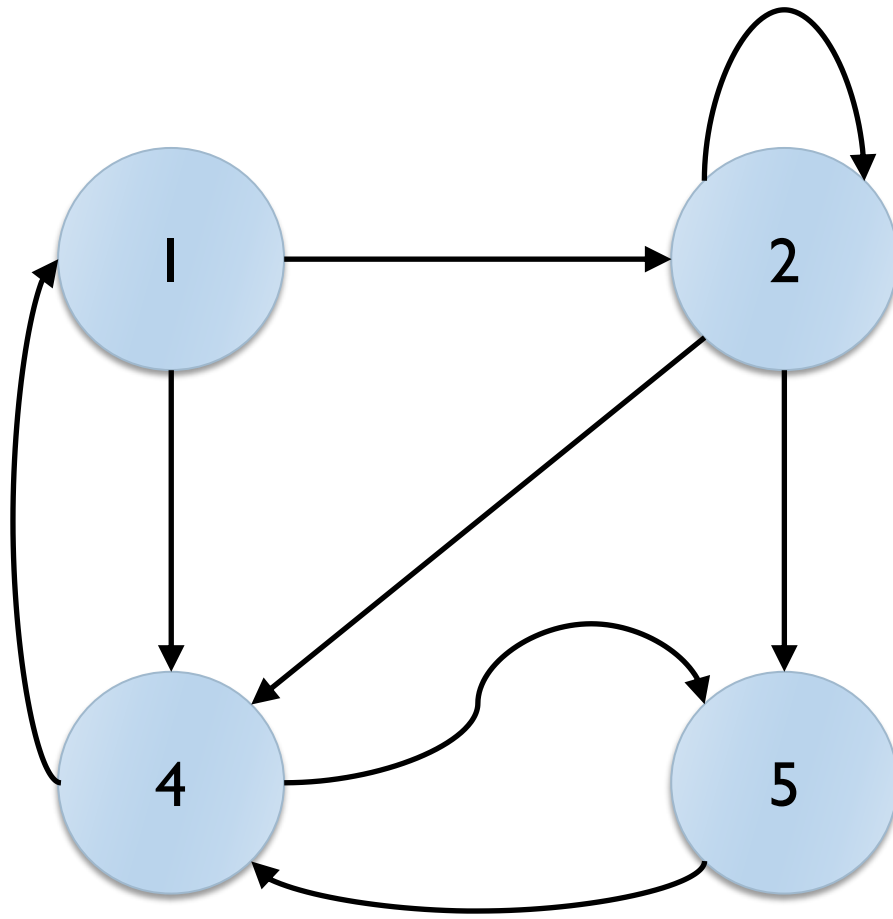
---

- ▶ A directed graph is **strongly connected** if, for every ordered pair of vertices  $(v, v')$ , there exists at least one path connecting  $v$  to  $v'$

# Example

---

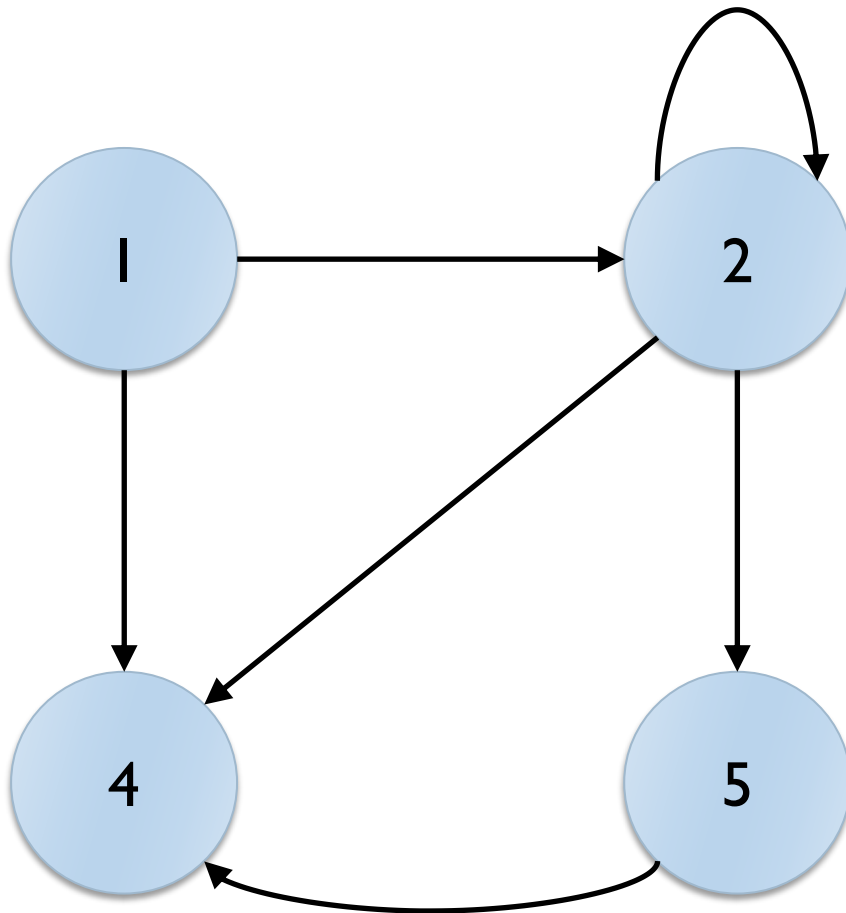
The graph is **strongly connected**



# Example

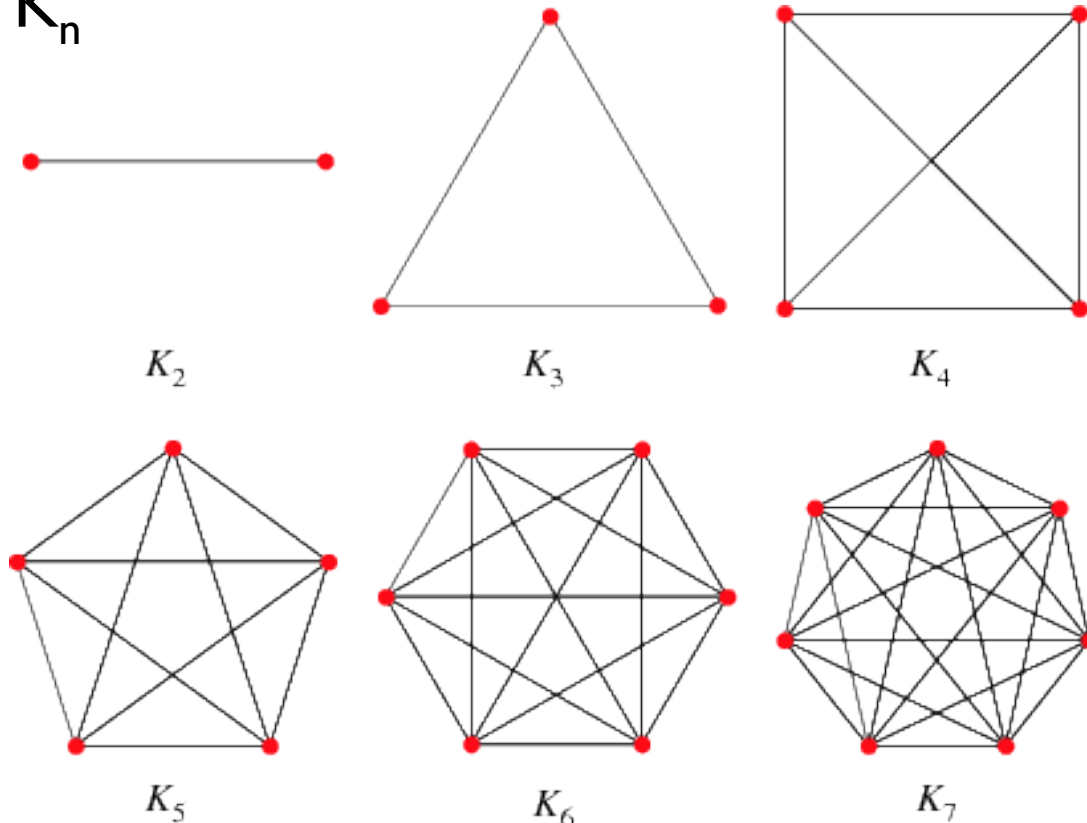
---

The graph is **not** strongly connected



# Complete graph

- ▶ A graph is complete if, for every pair of vertices, there is an edge connecting them (they are adjacent)
- ▶ Symbol:  $K_n$



# Complete graph: edges

---

- ▶ In a **complete** graph with  $n$  vertices, the number of **edges** is

	Directed	Undirected
No self loops	$n(n - 1)$	$\frac{n(n - 1)}{2}$
With self loops	$n^2$	$\frac{n(n + 1)}{2}$



# Density

---

- ▶ The density of a graph  $G=(V,E)$  is the ratio of the number of edges to the total number of possible edges

$$d = \frac{|E(G)|}{|E(K_{|V(G)|})|}$$

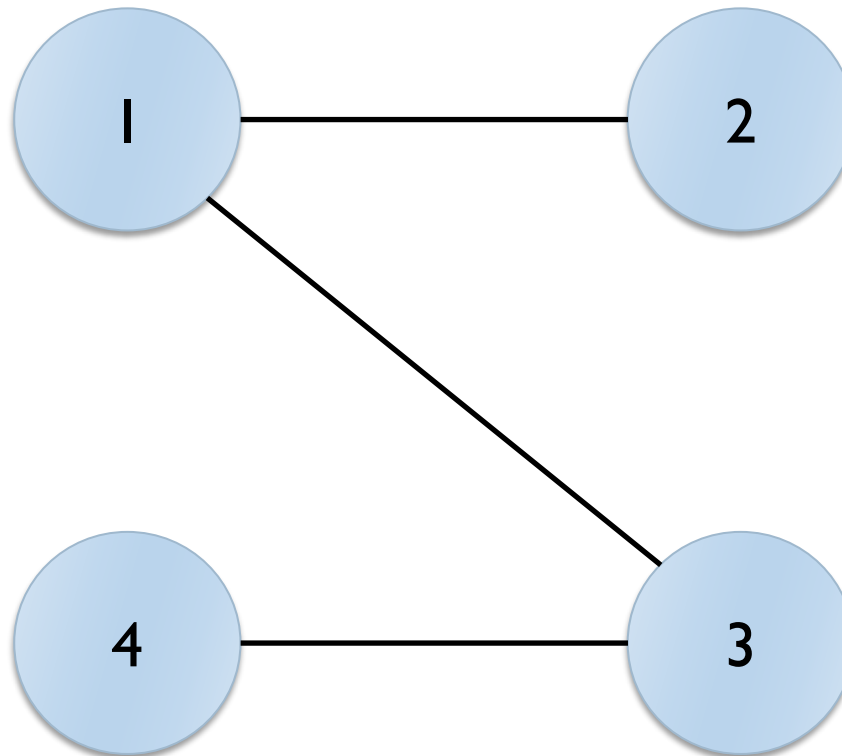
# Example

---

Density = 0.5

Existing: 3 edges

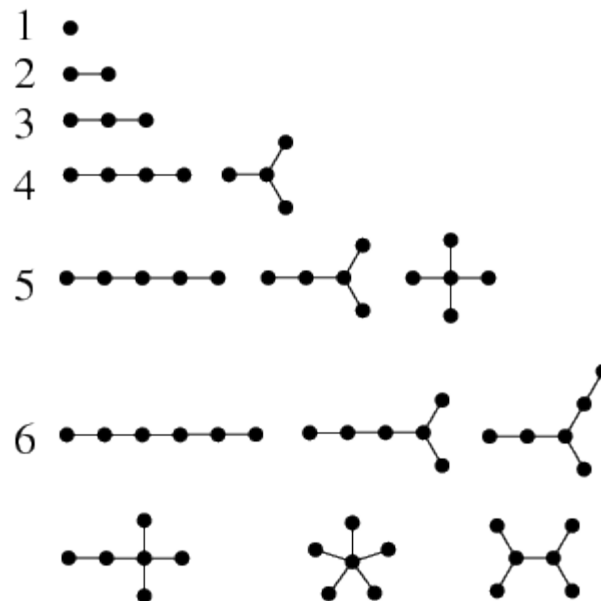
Total: 6 possible edges



# Trees and Forests

---

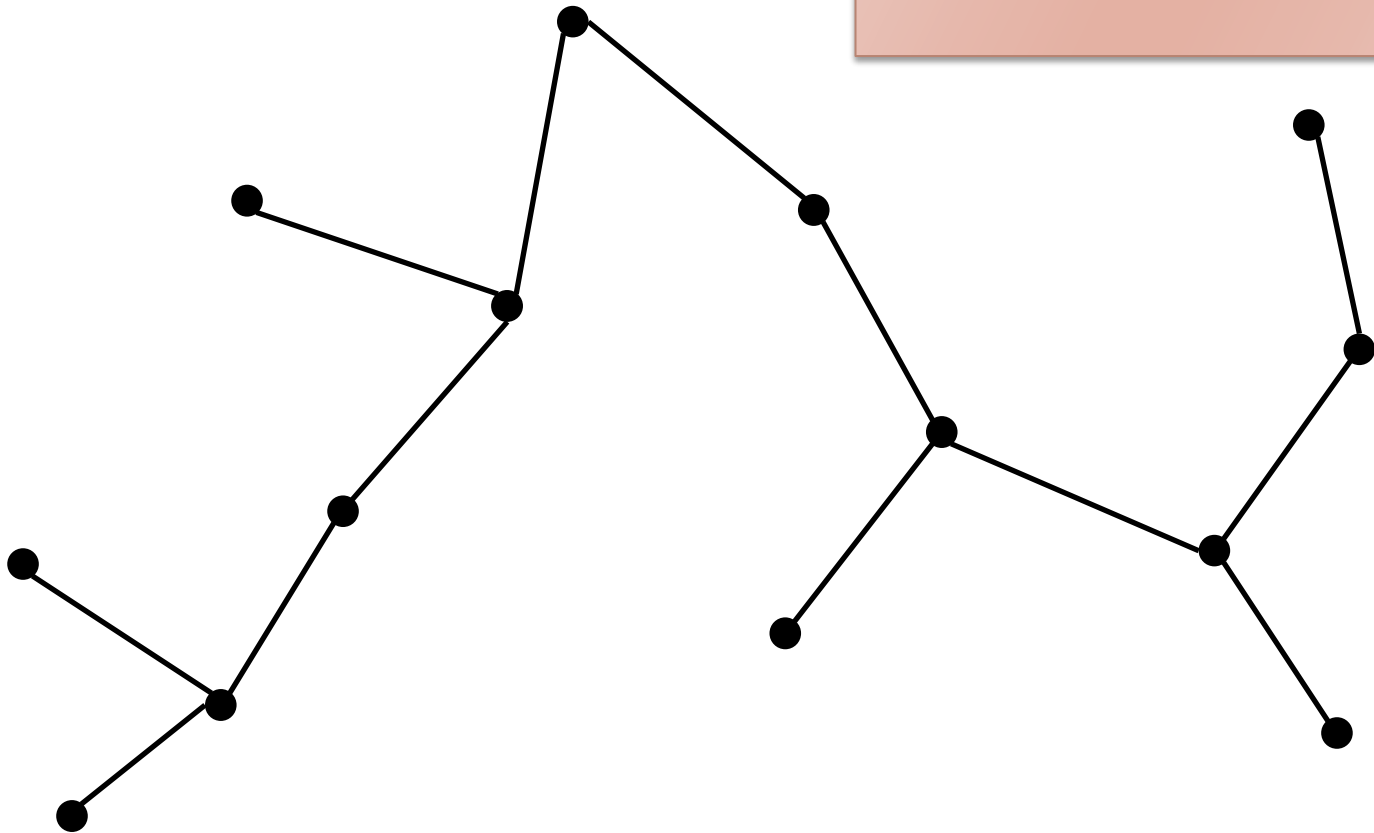
- ▶ An undirected acyclic graph is called **forest**
- ▶ An undirected acyclic connected graph is called **tree**



# Example

---

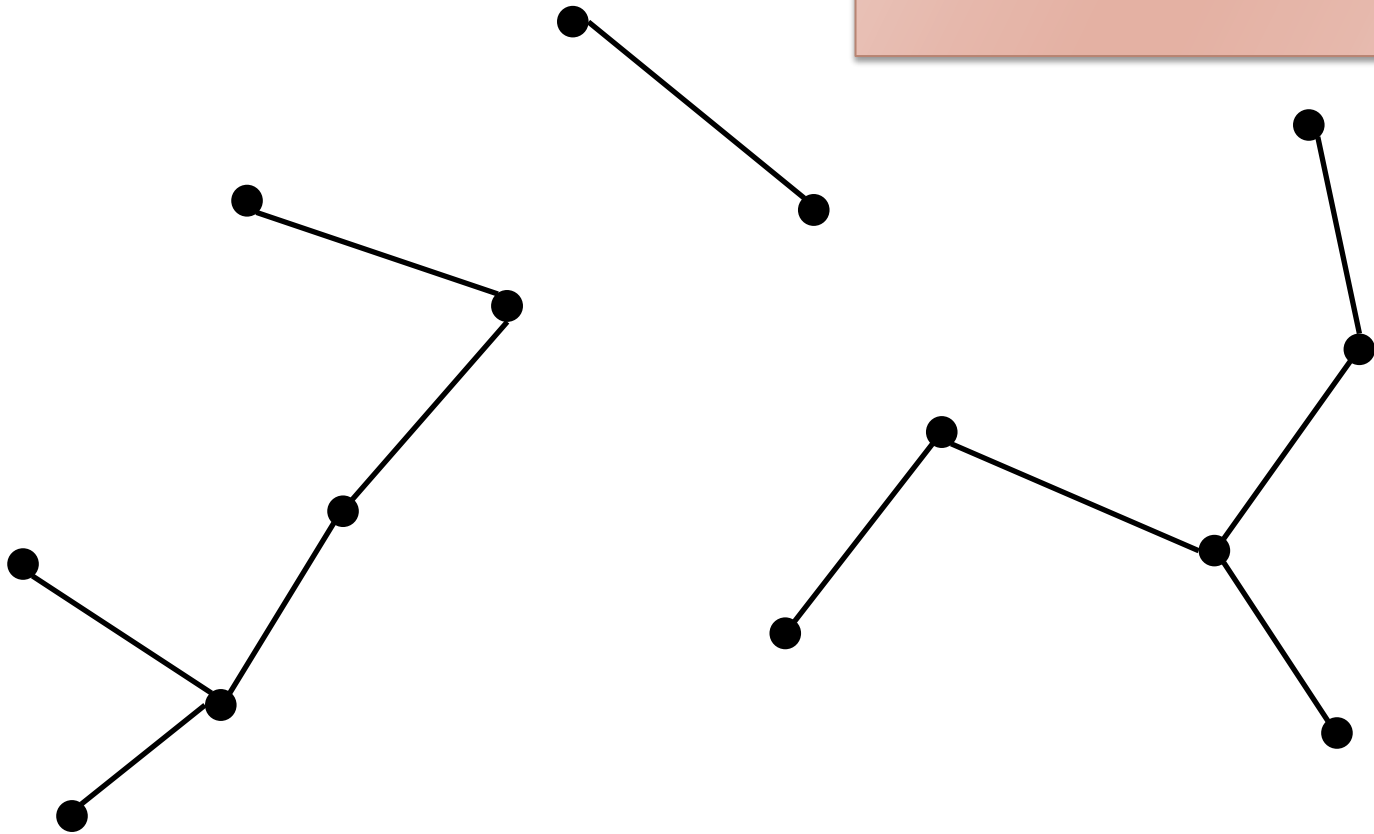
Tree



# Example

---

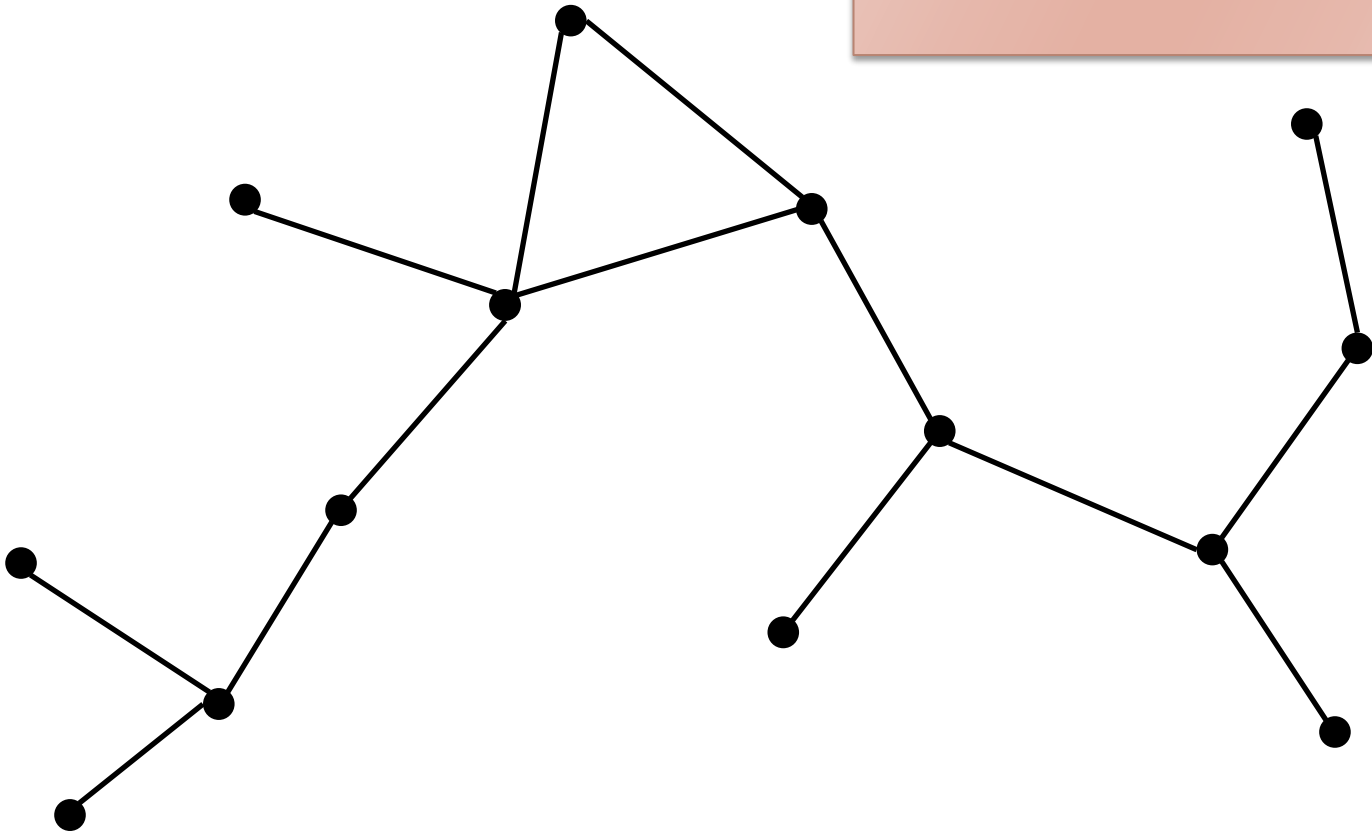
Forest



# Example

---

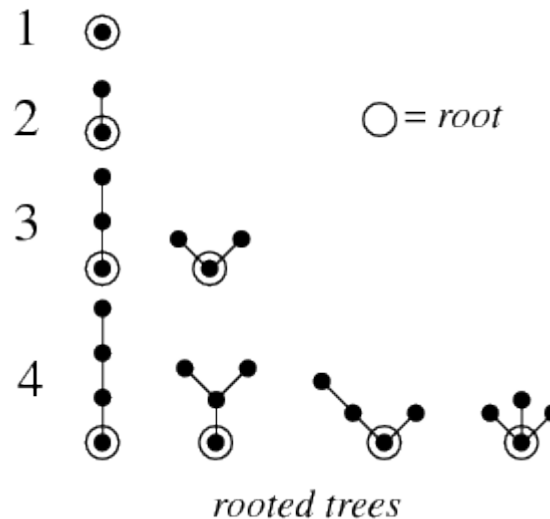
This is not a tree nor a forest  
(it contains a cycle)



# Rooted trees

---

- ▶ In a tree, a special node may be singled out
- ▶ This node is called the “**root**” of the tree
- ▶ Any node of a tree can be the root



# Tree (implicit) ordering

---

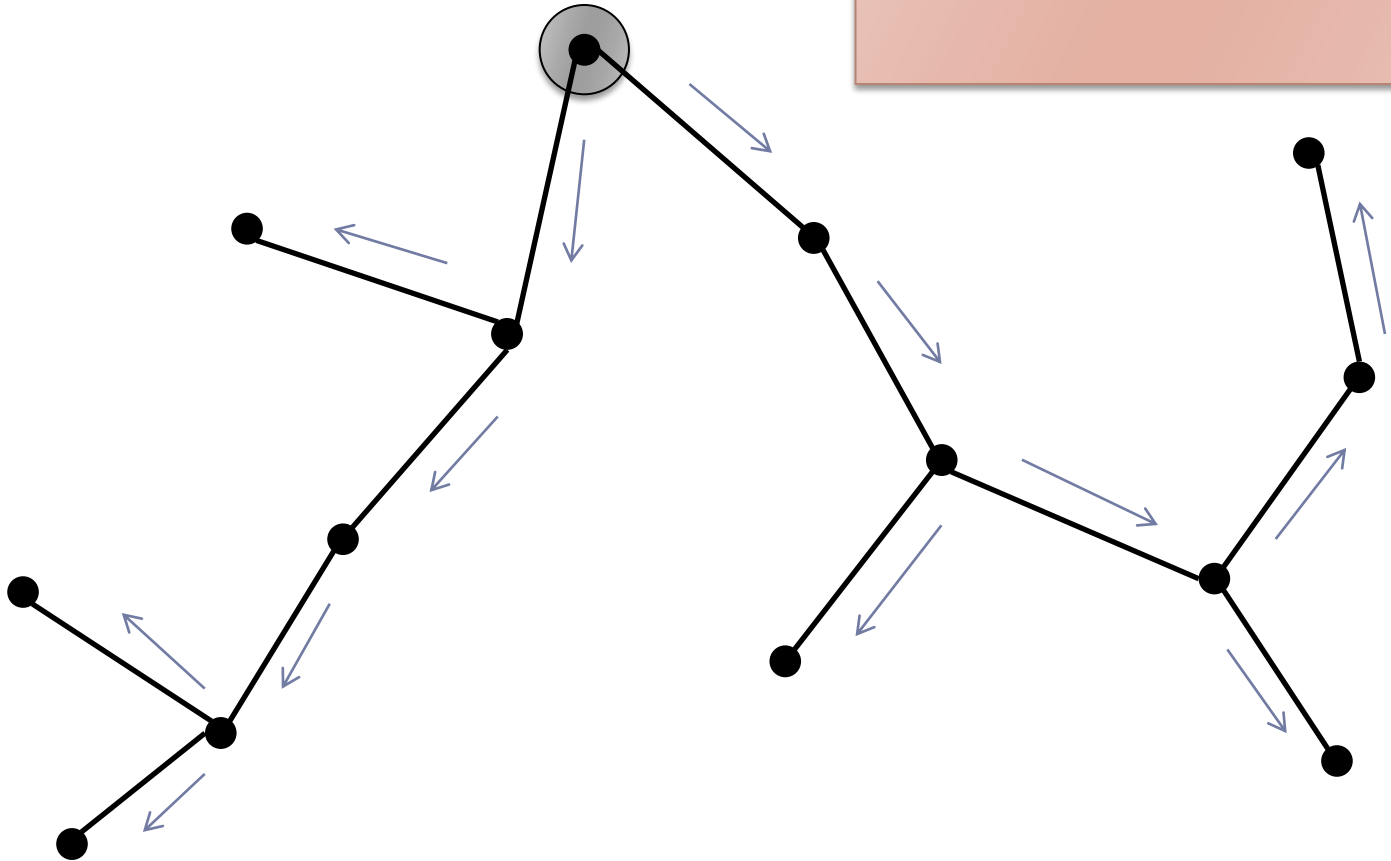
- ▶ The root node of a tree **induces an ordering** of the nodes
- ▶ The root is the “ancestor” of all other nodes/vertices
  - ▶ “children” are “away from the root”
  - ▶ “parents” are “towards the root”
- ▶ The root is the only node without parents
- ▶ All other nodes have exactly one parent
- ▶ The furthestmost (children-of-children-of-children...) nodes are “leaves”



# Example

---

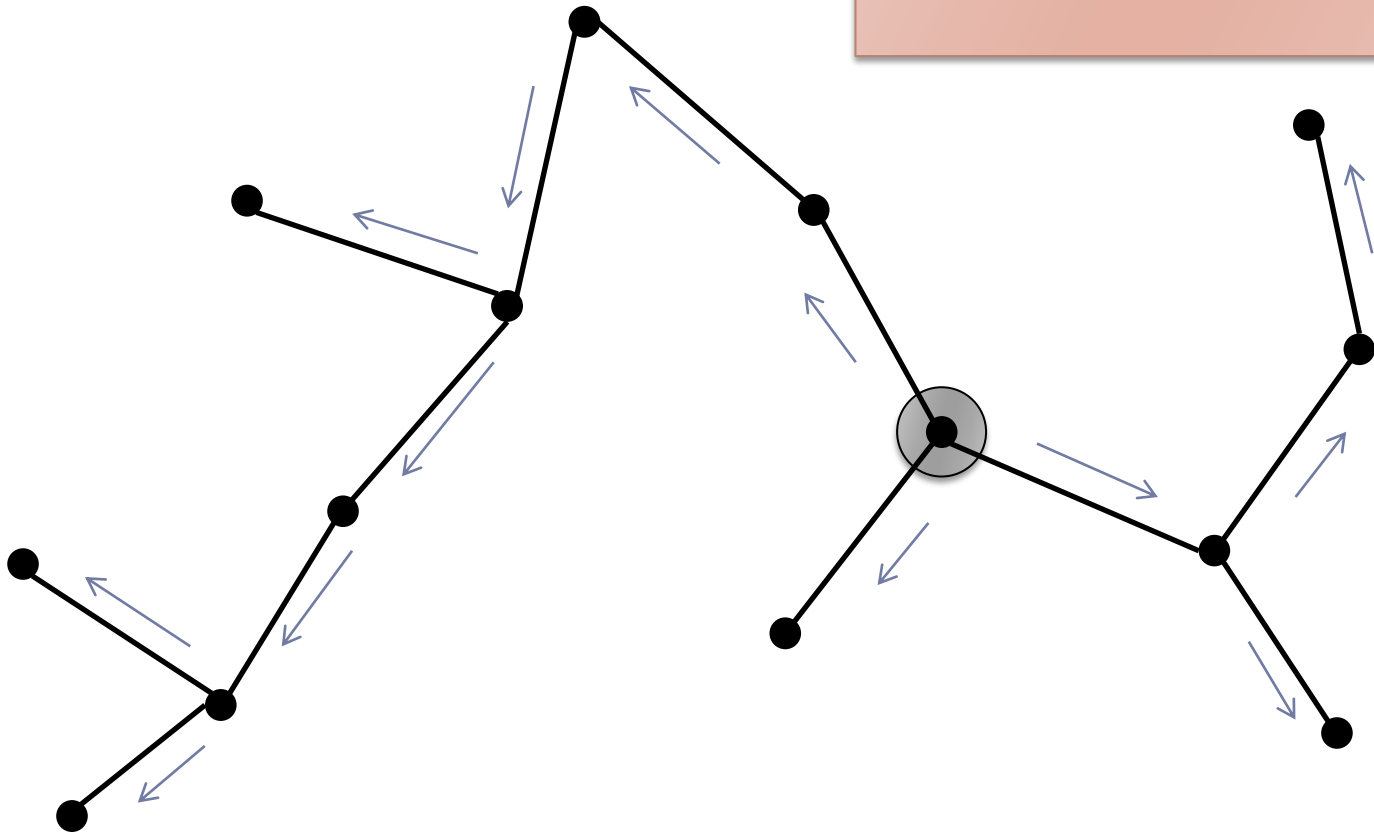
Rooted Tree



# Example

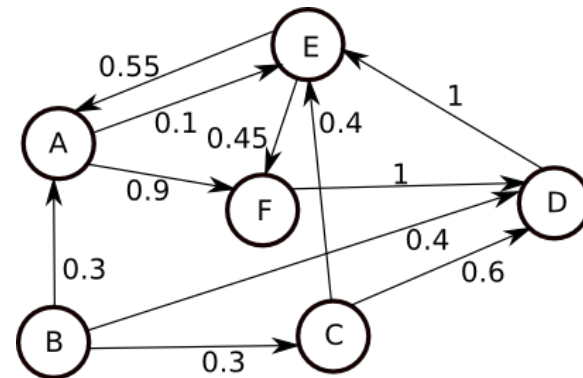
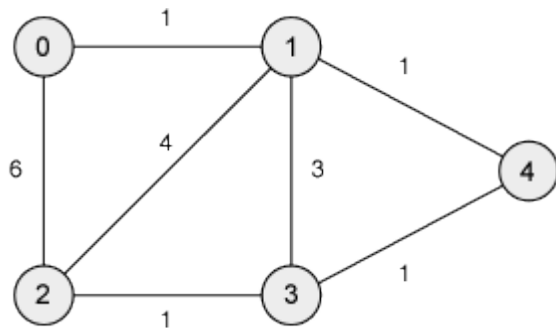
---

Rooted Tree








# Weighted graphs

- ▶ A weighted graph is a graph in which each branch (edge) is given a numerical weight.
- ▶ A weighted graph is therefore a special type of labeled graph in which the labels are numbers (which are usually taken to be positive).



# Licenza d'uso



- ▶ Queste diapositive sono distribuite con licenza Creative Commons “Attribuzione - Non commerciale - Condividi allo stesso modo (CC BY-NC-SA)”
- ▶ Sei libero:
  - ▶ di riprodurre, distribuire, comunicare al pubblico, esporre in pubblico, rappresentare, eseguire e recitare quest'opera 
  - ▶ di modificare quest'opera 
- ▶ Alle seguenti condizioni:
  - ▶ **Attribuzione** — Devi attribuire la paternità dell'opera agli autori originali e in modo tale da non suggerire che essi avallino te o il modo in cui tu usi l'opera. 
  - ▶ **Non commerciale** — Non puoi usare quest'opera per fini commerciali. 
  - ▶ **Condividi allo stesso modo** — Se alteri o trasformi quest'opera, o se la usi per crearne un'altra, puoi distribuire l'opera risultante solo con una licenza identica o equivalente a questa. 
- ▶ <http://creativecommons.org/licenses/by-nc-sa/3.0/>