



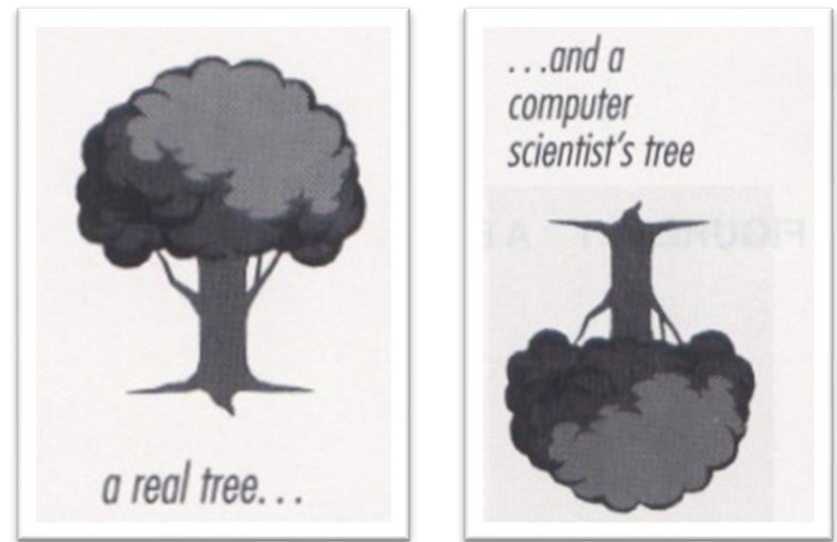
Trees

(in computer science)



Tree in Computer Science

- ▶ A tree is a widely used data structure that simulates a hierarchical tree structure with a set of linked nodes

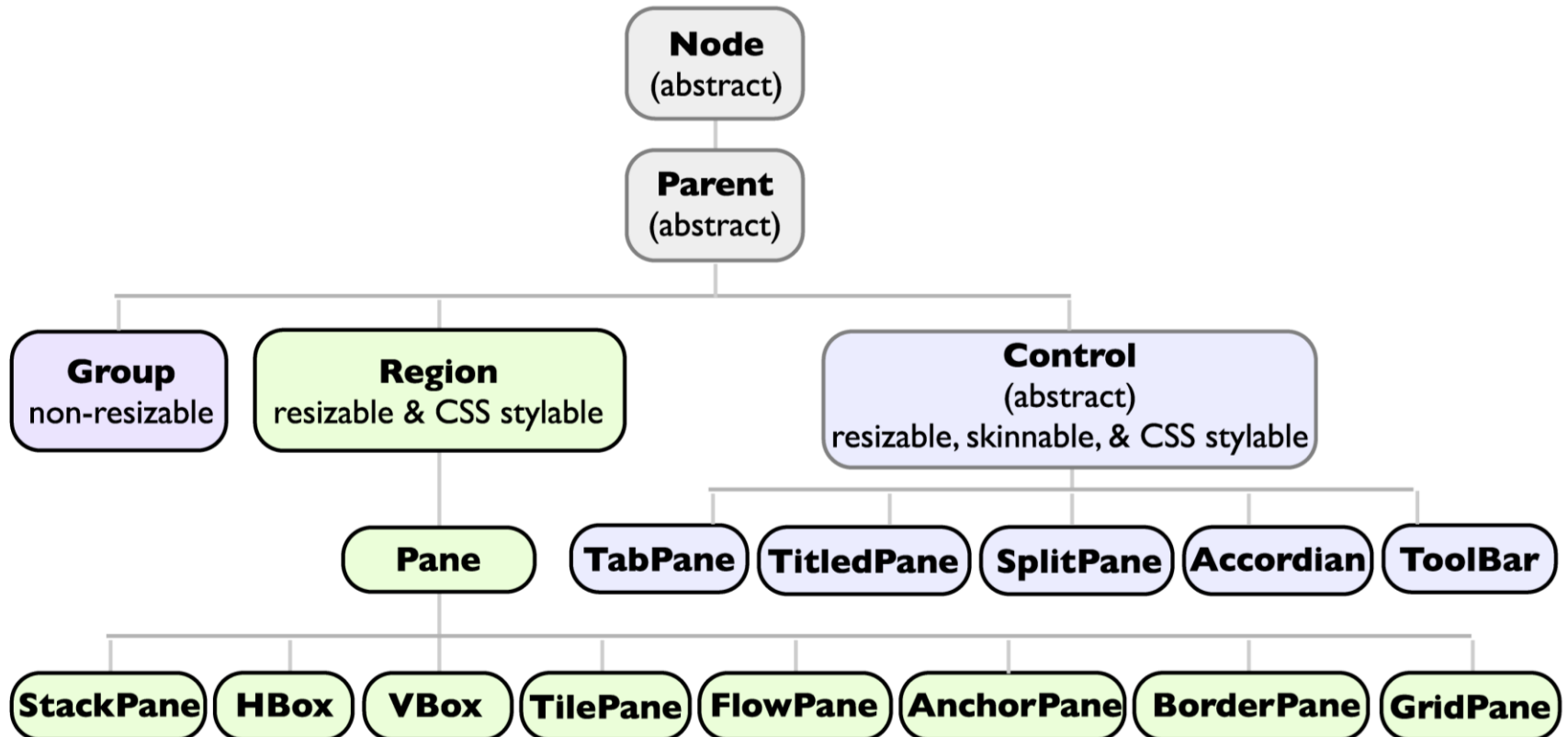


Tree in Computer Science

- ▶ **Fundamental** data storage structures used in programming
- ▶ Nonlinear structure
- ▶ Represents a *hierarchy*
- ▶ Items in a tree do not form a simple sequence
- ▶ Quite efficient for **retrieving** items (as arrays)
- ▶ Quite efficient for **inserting/deleting** items (as lists)



JavaFX 2.0 Layout Classes



Ordinamento dello Stato Italiano



Tree basics

- ▶ Consists of nodes connected by edges
- ▶ Nodes often represent entities (complex objects)
- ▶ Edges between the nodes represent the way the nodes are related
- ▶ The only way to get from node to node is to follow a path along the edges

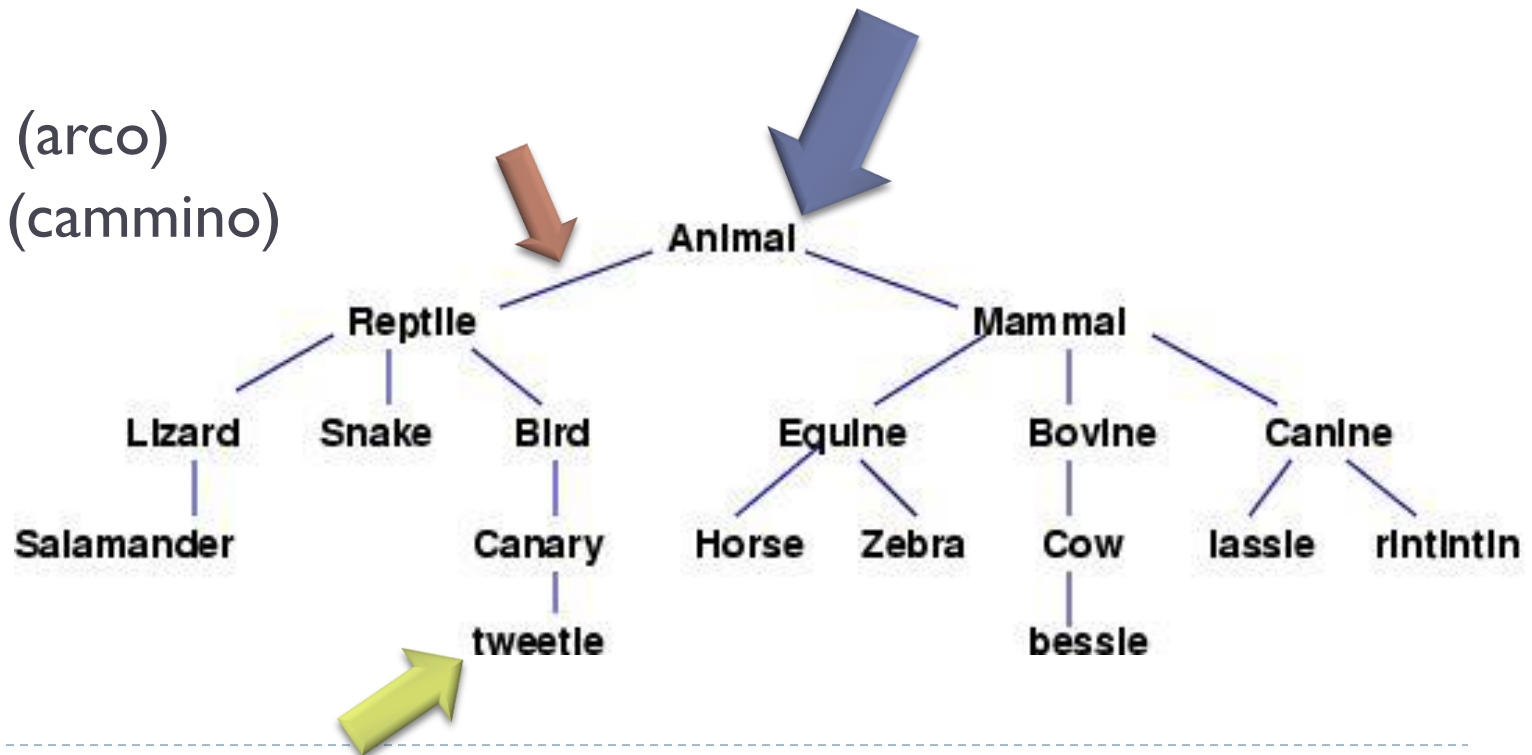
Tree Basics

▶ Node

- ▶ Root (radice)
- ▶ Leaf (foglia)
- ▶ Interior node/branch (nodo interno)

▶ Links

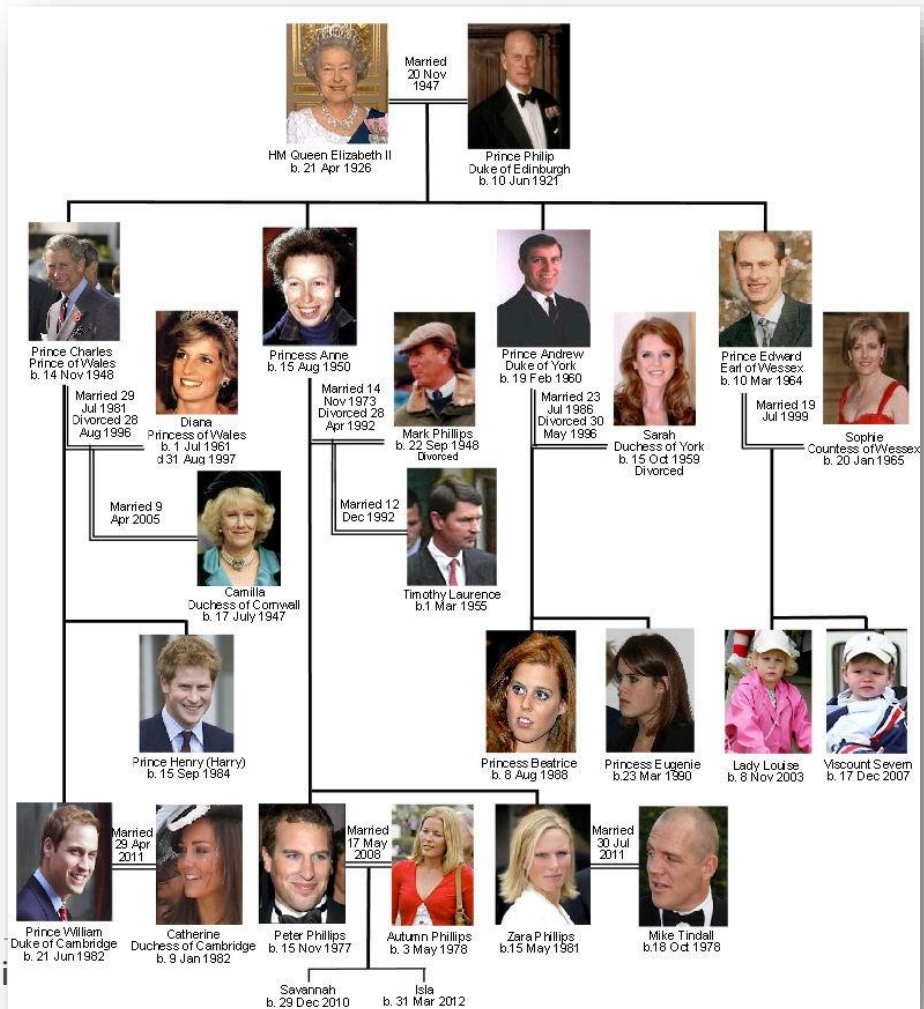
- ▶ Edge (arco)
- ▶ Path (cammino)

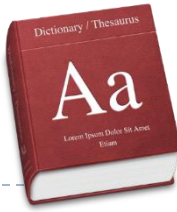


Tree Basics

► Relationship

- Parent (padre)
- Child nodes (nodi figli)
- Sibling (fratelli)
- Descendant (discendente, successore)
- Ancestor (antenato, predecessore)





Terminology

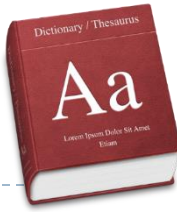
▶ Visiting

- ▶ A node is visited when program control arrives at the node, usually for processing

▶ Traversing

- ▶ To traverse a tree means to visit all the nodes in some specified order

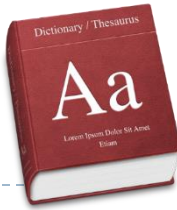
Terminology



► Levels

- The level of a particular node refers to how many generations the node is from the root
- Root is assumed to be level 0

Terminology



► Height

- The height of a node is the length of the path to its farthest descendant (i.e. farthest leaf node)
- The height of a tree is the height of the root
- A tree with only root node has height 0

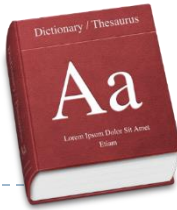


Binary Trees

Binary Tree

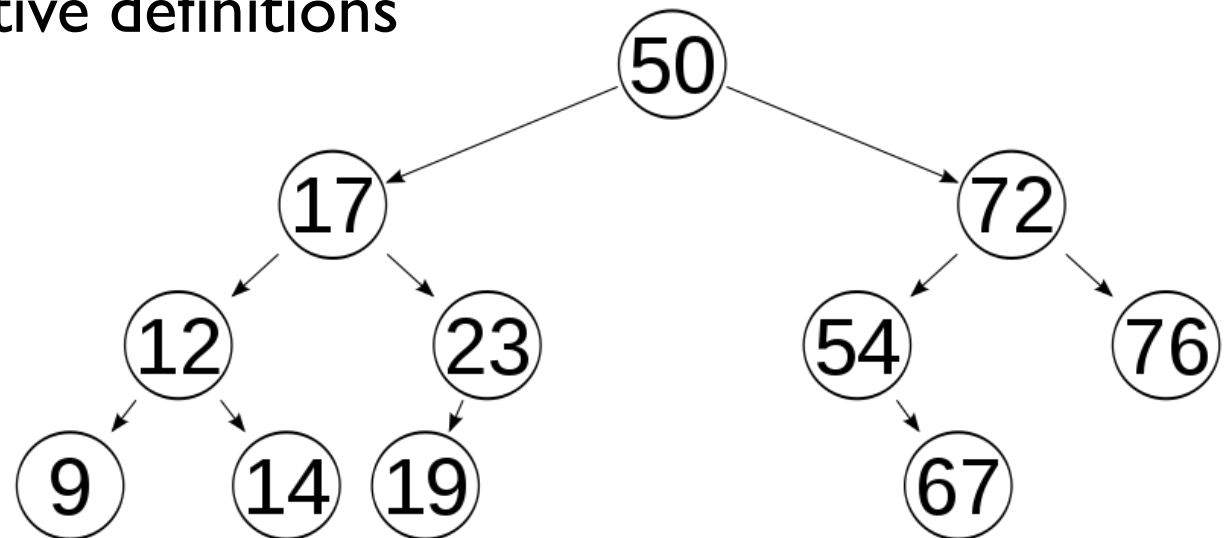
- ▶ A binary tree is a tree where each node has at most two children
- ▶ The two children are ordered (“left”, “right”)
 - ▶ Right sub-tree vs. Left sub-tree

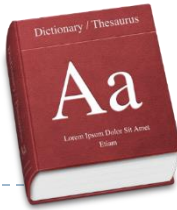




Balanced trees

- ▶ (Height-)balanced trees
 - ▶ The left and right sub-trees' heights differ by at most one
 - ▶ The two sub-trees are (height-)balanced
- ▶ Perfectly balanced
 - ▶ $2^h - 1$ nodes
- ▶ Several alternative definitions

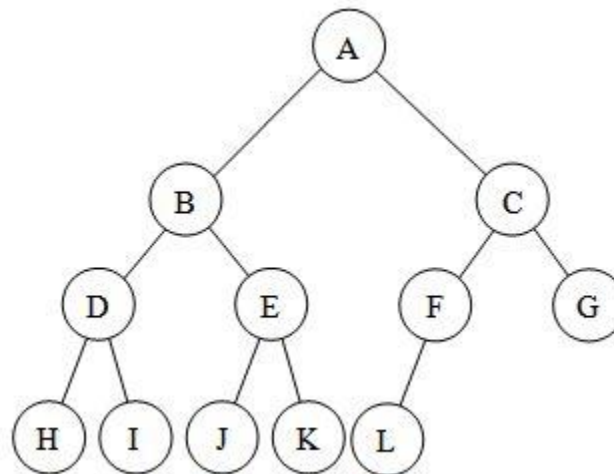




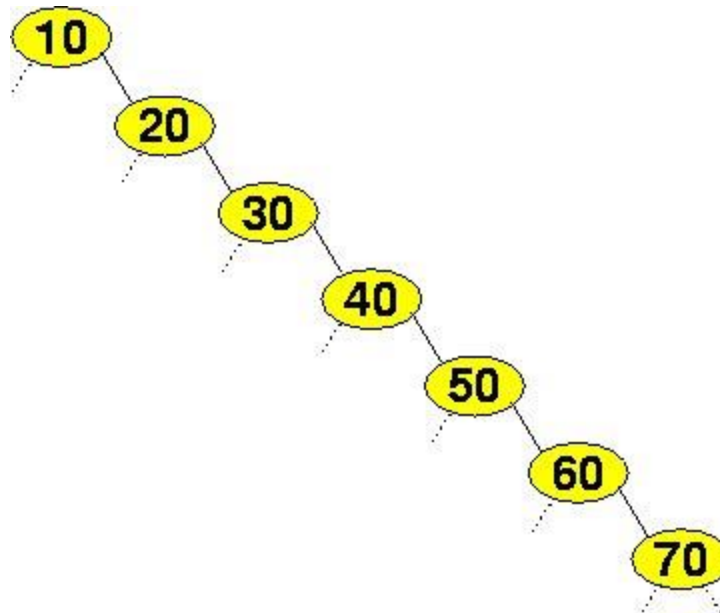
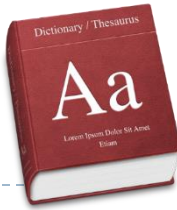
Complete trees

► Complete binary tree

- Every level, except possibly the last, is completely filled, and all nodes are as far left as possible



Degenerate trees



Traversal in binary trees

- ▶ **Pre-order**

- ▶ process root node, then its left/right sub-trees

- ▶ **In-order**

- ▶ process left sub-tree, then root node, then right

- ▶ **Post-order**

- ▶ process left/right sub-trees, then root node



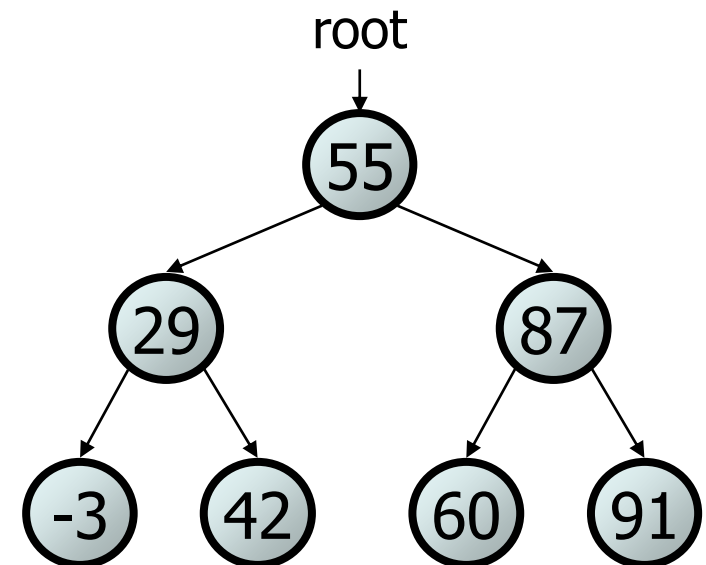


BST

Binary Search Tree

Binary search trees

- ▶ A binary tree where each non-empty node R has the following properties:
 - ▶ Elements of R 's left sub-tree contain data “less than” R 's data
 - ▶ Elements of R 's right sub-tree contain data “greater than” R 's
 - ▶ R 's left and right sub-trees are also binary search trees

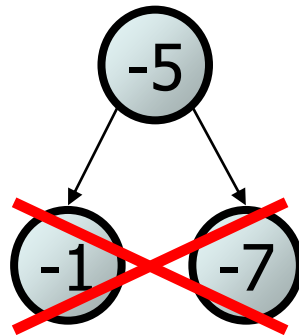


Binary search trees

- ▶ BSTs store their elements in sorted order, which is helpful for searching/sorting tasks

Exercise

- Is it a legal binary search tree?



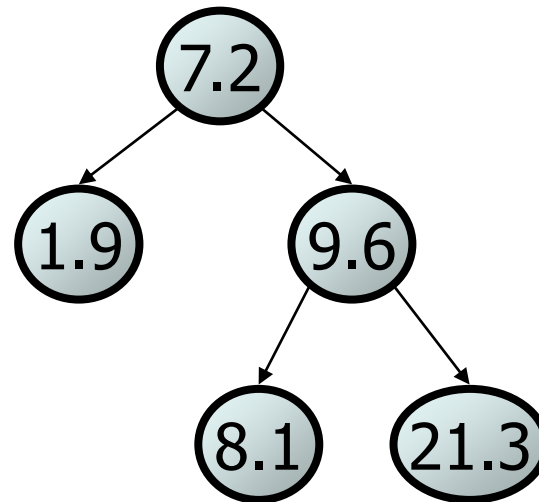
Exercise

- ▶ Is it a legal binary search tree?

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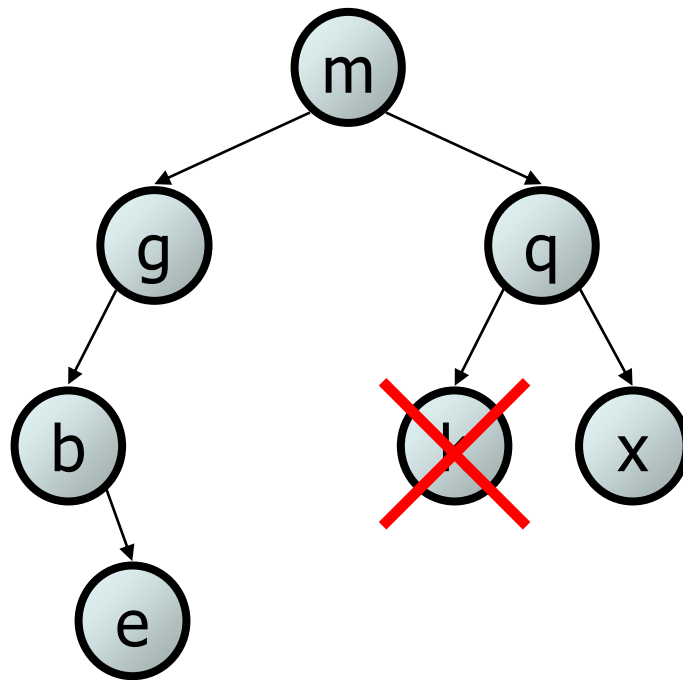
Exercise

- Is it a legal binary search tree?



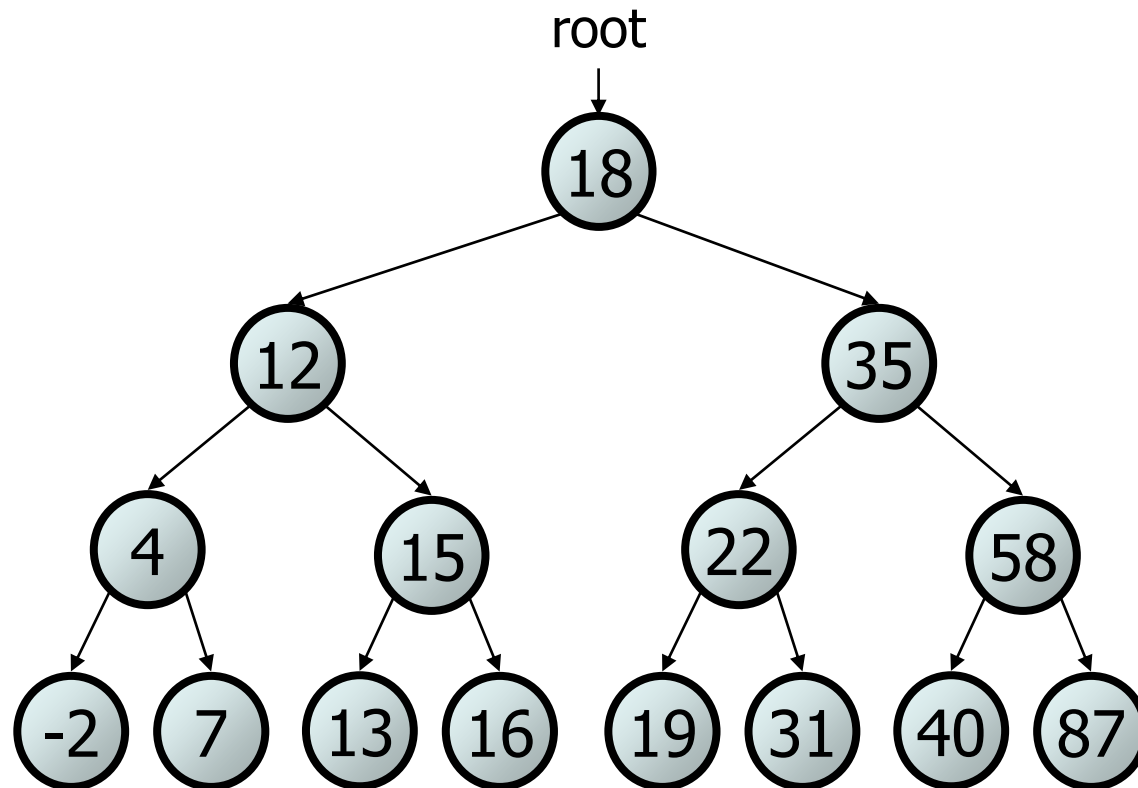
Exercise

- Is it a legal binary search tree?



Searching in a BST

- Describe an algorithm for searching a binary search tree (try searching for 31, then 6)



Searching in a BST

- ▶ Searching in a BST is $O(h)$

If the tree is balanced, then $h \cong \log_2 N$

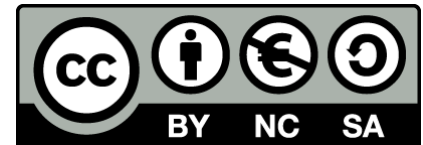
\Rightarrow Searching for an element is $O(\ln N)$








Showdown

	Array	List	Hash	BST
add(element)	$O(1)$	$O(1)$	$O(1)$	$O(\ln n)$
remove(object)	$O(n) + O(n)$	$O(n) + O(1)$	$O(1)$	$O(\ln n)$
get(index)	$O(1)$	$O(n)$	n.a.	n.a.
set(index, element)	$O(1)$	$O(n) + O(1)$	n.a.	n.a.
add(index, element)	$O(1) + O(n)$	$O(n) + O(1)$	n.a.	n.a.
remove(index)	$O(n)$	$O(n) + O(1)$	n.a.	n.a.
contains(object)	$O(n)$	$O(n)$	$O(1)$	$O(\ln n)$
indexOf(object)	$O(n)$	$O(n)$	n.a.	n.a.

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